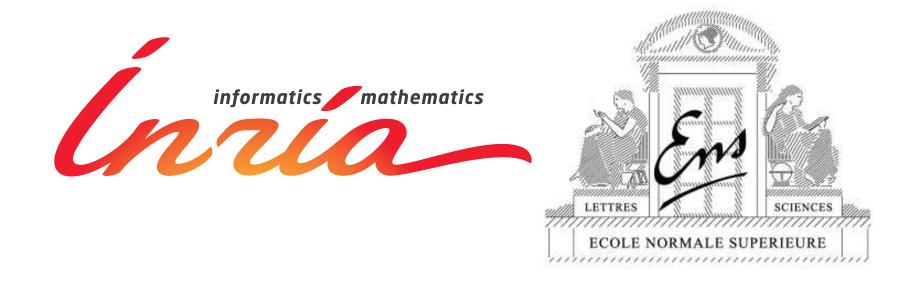
Machine learning and convex optimization with submodular functions

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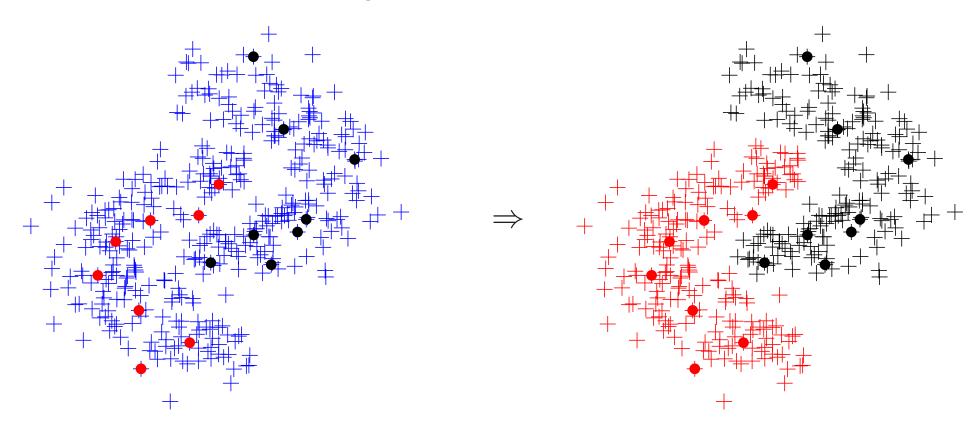
Workshop on combinatorial optimization - Cargese, 2013

Submodular functions - References

- References based on combinatorial optimization
 - Submodular Functions and Optimization (Fujishige, 2005)
 - Discrete convex analysis (Murota, 2003)
- Tutorial paper based on convex optimization (Bach, 2011b)
 - www.di.ens.fr/~fbach/submodular_fot.pdf
- Slides for this lecture
 - www.di.ens.fr/~fbach/fbach_cargese_2013.pdf

Submodularity (almost) everywhere Clustering

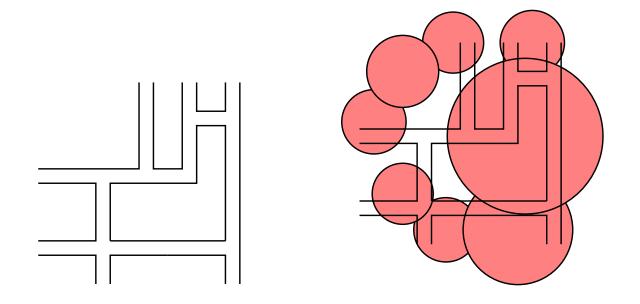
• Semi-supervised clustering



• Submodular function minimization

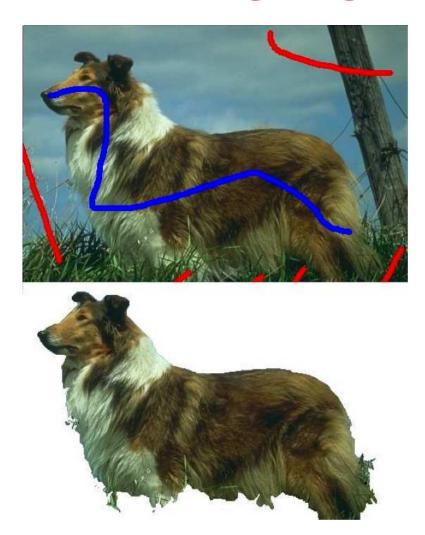
Submodularity (almost) everywhere Sensor placement

- Each sensor covers a certain area (Krause and Guestrin, 2005)
 - Goal: maximize coverage



- Submodular function maximization
- Extension to experimental design (Seeger, 2009)

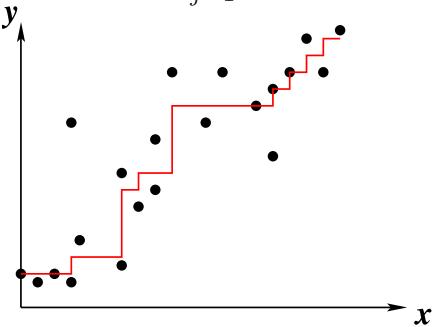
Submodularity (almost) everywhere Graph cuts and image segmentation



Submodular function minimization

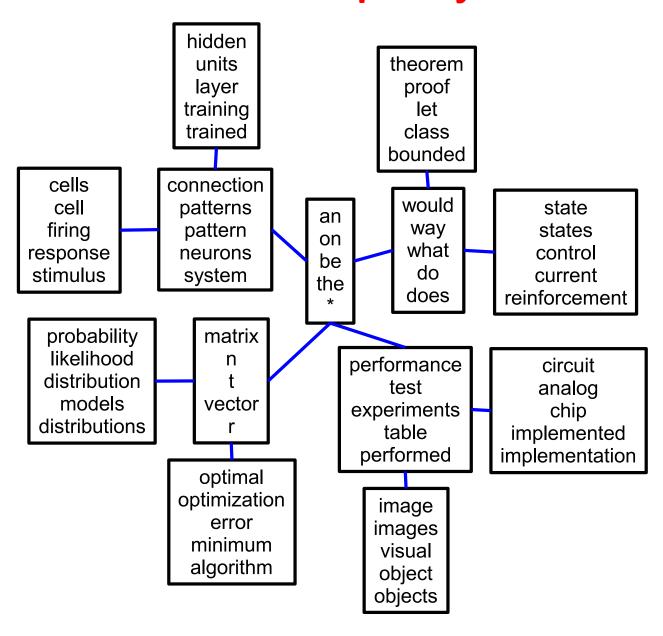
Submodularity (almost) everywhere Isotonic regression

- Given real numbers x_i , $i = 1, \ldots, p$
 - Find $y \in \mathbb{R}^p$ that minimizes $\frac{1}{2} \sum_{j=1}^p (x_i y_i)^2$ such that $\forall i, y_i \leqslant y_{i+1}$

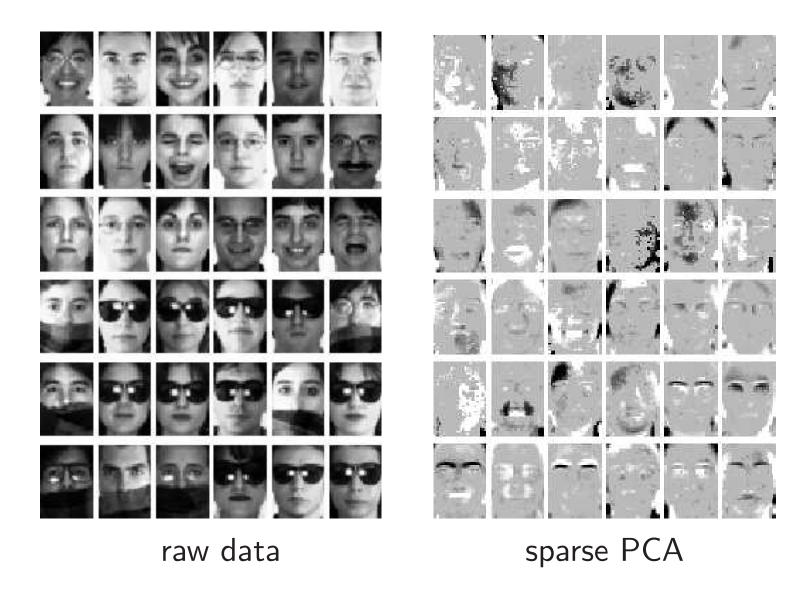


Submodular convex optimization problem

Submodularity (almost) everywhere Structured sparsity - I

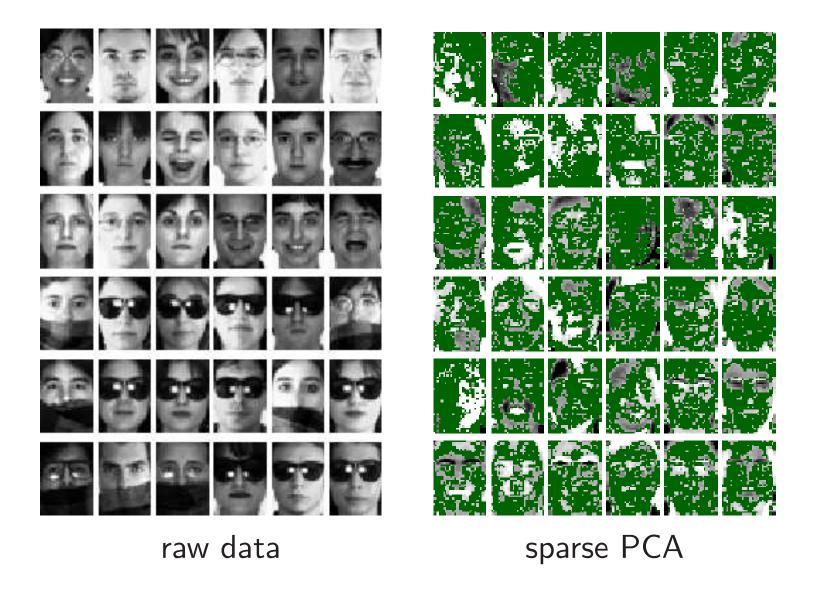


Submodularity (almost) everywhere Structured sparsity - II



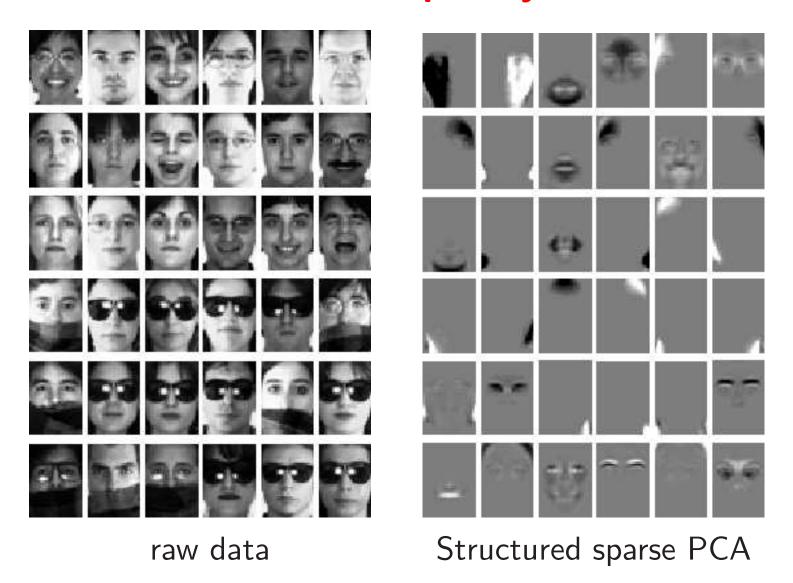
No structure: many zeros do not lead to better interpretability

Submodularity (almost) everywhere Structured sparsity - II



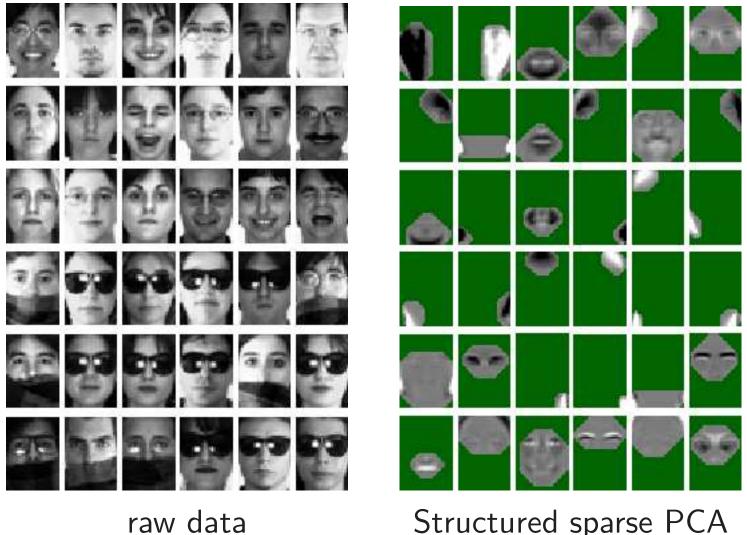
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Submodularity (almost) everywhere Structured sparsity - II



• Submodular convex optimization problem

Submodularity (almost) everywhere Structured sparsity - II



Structured sparse PCA

Submodular convex optimization problem

Submodularity (almost) everywhere Image denoising

• Total variation denoising (Chambolle, 2005)

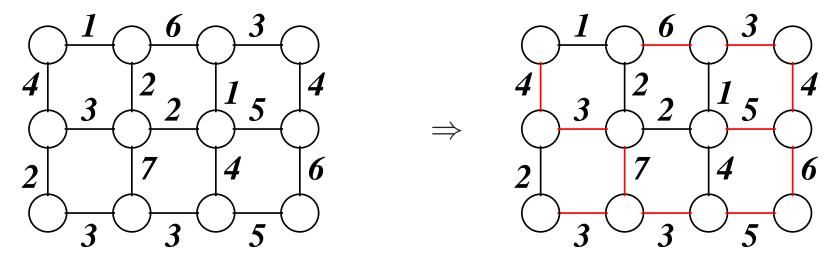




• Submodular convex optimization problem

Submodularity (almost) everywhere Maximum weight spanning trees

- Given an undirected graph G = (V, E) and weights $w : E \mapsto \mathbb{R}_+$
 - find the maximum weight spanning tree



Greedy algorithm for submodular polyhedron - matroid

Submodularity (almost) everywhere Combinatorial optimization problems

- Set $V = \{1, ..., p\}$
- ullet Power set $2^V=$ set of all subsets, of cardinality 2^p
- Minimization/maximization of a set function $F: 2^V \to \mathbb{R}$.

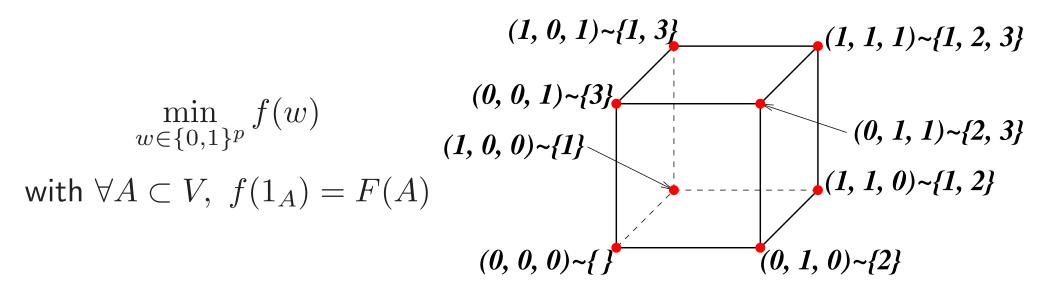
$$\min_{A \subset V} F(A) = \min_{A \in 2^V} F(A)$$

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Reformulation as (pseudo) Boolean function



Submodularity (almost) everywhere Convex optimization with combinatorial structure

- Supervised learning / signal processing
 - Minimize regularized empirical risk from data (x_i, y_i) , $i = 1, \ldots, n$:

$$\min_{f \in \mathcal{F}} \frac{1}{n} \sum_{i=1}^{n} \ell(y_i, f(x_i)) + \lambda \Omega(f)$$

- $-\mathcal{F}$ is often a vector space, formulation often convex
- Introducing discrete structures within a vector space framework
 - Trees, graphs, etc.
 - Many different approaches (e.g., stochastic processes)
- Submodularity allows the incorporation of discrete structures

Outline

1. Submodular functions

- Review and examples of submodular functions
- Links with convexity through Lovász extension

2. Submodular minimization

- Non-smooth convex optimization
- Parallel algorithm for special case

3. Structured sparsity-inducing norms

- Relaxation of the penalization of supports by submodular functions
- Extensions (symmetric, ℓ_q -relaxation)

Submodular functions Definitions

• **Definition**: $F: 2^V \to \mathbb{R}$ is **submodular** if and only if

$$\forall A, B \subset V, \quad F(A) + F(B) \geqslant F(A \cap B) + F(A \cup B)$$

- NB: equality for modular functions
- Always assume $F(\varnothing) = 0$

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• Equivalent definition:

 $\forall k \in V, \quad A \mapsto F(A \cup \{k\}) - F(A)$ is non-increasing

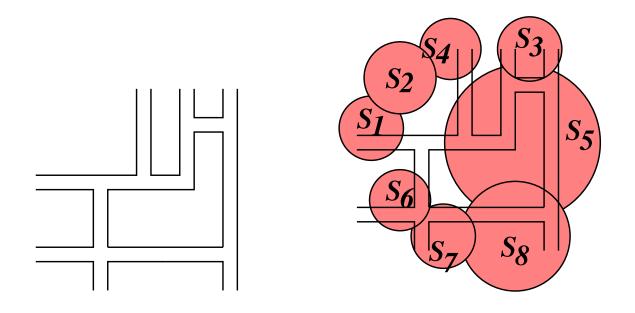
$$\Leftrightarrow \forall A \subset B, \ \forall k \notin A, \ F(A \cup \{k\}) - F(A) \geqslant F(B \cup \{k\}) - F(B)$$

"Concave property": Diminishing return property

Examples of submodular functionsCardinality-based functions

- Notation for modular function: $s(A) = \sum_{k \in A} s_k$ for $s \in \mathbb{R}^p$
 - If $s = 1_V$, then s(A) = |A| (cardinality)
- **Proposition**: If $s \in \mathbb{R}^p_+$ and $g : \mathbb{R}_+ \to \mathbb{R}$ is a concave function, then $F : A \mapsto g(s(A))$ is submodular
- **Proposition 2**: If $F:A\mapsto g(s(A))$ is submodular for all $s\in\mathbb{R}^p_+$, then g is concave
- Classical example:
 - -F(A)=1 if |A|>0 and 0 otherwise
 - May be rewritten as $F(A) = \max_{k \in V} (1_A)_k$

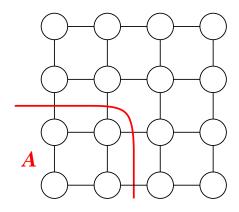
Examples of submodular functions Covers



- Let W be any "base" set, and for each $k \in V$, a set $S_k \subset W$
- Set cover defined as $F(A) = |\bigcup_{k \in A} S_k|$
- Proof of submodularity ⇒ homework

Examples of submodular functions Cuts

- ullet Given a (un)directed graph, with vertex set V and edge set E
 - -F(A) is the total number of edges going from A to $V \setminus A$.



• Generalization with $d: V \times V \to \mathbb{R}_+$

$$F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j)$$

Proof of submodularity ⇒ homework

Examples of submodular functions Entropies

- ullet Given p random variables X_1,\ldots,X_p with finite number of values
 - Define F(A) as the joint entropy of the variables $(X_k)_{k\in A}$
 - − F is submodular
- Proof of submodularity using data processing inequality (Cover and Thomas, 1991): if $A \subset B$ and $k \notin B$,

$$F(A \cup \{k\}) - F(A) = H(X_A, X_k) - H(X_A) = H(X_k | X_A) \ge H(X_k | X_B)$$

- Symmetrized version $G(A)=F(A)+F(V\backslash A)-F(V)$ is mutual information between X_A and $X_{V\backslash A}$
- Extension to continuous random variables, e.g., Gaussian: $F(A) = \log \det \Sigma_{AA}$, for some positive definite matrix $\Sigma \in \mathbb{R}^{p \times p}$

Examples of submodular functions Flows

- Net-flows from multi-sink multi-source networks (Megiddo, 1974)
- See details in Fujishige (2005); Bach (2011b)
- Efficient formulation for set covers

Examples of submodular functions Matroids

- The pair (V, \mathcal{I}) is a matroid with \mathcal{I} its family of independent sets, iff:
- (a) $\varnothing \in \mathcal{I}$
- (b) $I_1 \subset I_2 \in \mathcal{I} \Rightarrow I_1 \in \mathcal{I}$
- (c) for all $I_1, I_2 \in \mathcal{I}$, $|I_1| < |I_2| \Rightarrow \exists k \in I_2 \backslash I_1, \ I_1 \cup \{k\} \in \mathcal{I}$
- Rank function of the matroid, defined as $F(A) = \max_{I \subset A, A \in \mathcal{I}} |I|$ is submodular (direct proof)

Graphic matroid

- V edge set of a certain graph G = (U, V)
- $-\mathcal{I}=$ set of subsets of edges which do not contain any cycle
- F(A) = |U| minus the number of connected components of the subgraph induced by A

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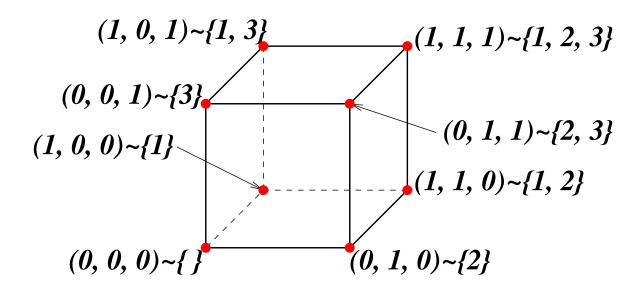
- Relaxation of the penalization of supports by submodular functions
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Choquet integral (Choquet, 1954) - Lovász extension

- Subsets may be identified with elements of $\{0,1\}^p$
- Given any set-function F and w such that $w_{j_1} \geqslant \cdots \geqslant w_{j_p}$, define:

$$f(w) = \sum_{k=1}^{p} w_{j_k} [F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})]$$

$$= \sum_{k=1}^{p-1} (w_{j_k} - w_{j_{k+1}}) F(\{j_1, \dots, j_k\}) + w_{j_p} F(\{j_1, \dots, j_p\})$$



Choquet integral (Choquet, 1954) - Lovász extension Properties

$$f(w) = \sum_{k=1}^{p} w_{j_k} [F(\{j_1, \dots, j_k\}) - F(\{j_1, \dots, j_{k-1}\})]$$

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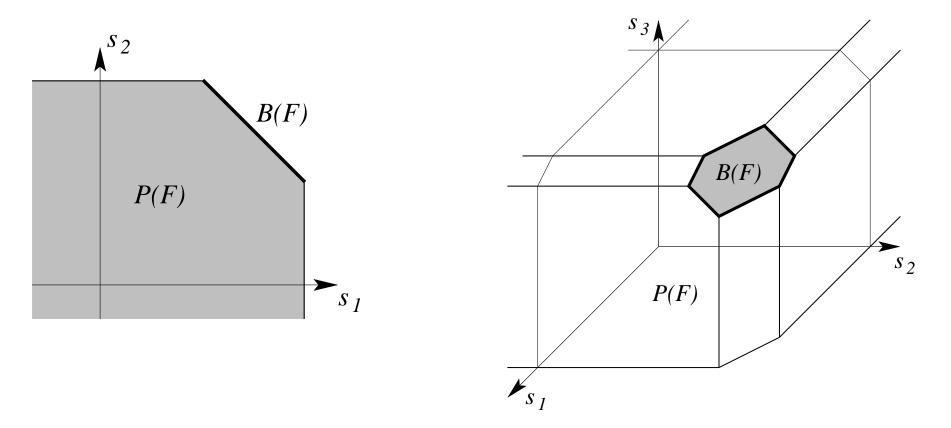
- ullet For any set-function F (even not submodular)
 - -f is piecewise-linear and positively homogeneous
 - If $w=1_A$, $f(w)=F(A)\Rightarrow$ extension from $\{0,1\}^p$ to \mathbb{R}^p

Submodular functions Links with convexity (Edmonds, 1970; Lovász, 1982)

- Theorem (Lovász, 1982): F is submodular if and only if f is convex
- Proof requires additional notions from Edmonds (1970):
 - Submodular and base polyhedra

Submodular and base polyhedra - Definitions

- Submodular polyhedron: $P(F) = \{s \in \mathbb{R}^p, \ \forall A \subset V, \ s(A) \leqslant F(A)\}$
- Base polyhedron: $B(F) = P(F) \cap \{s(V) = F(V)\}$



ullet Property: P(F) has non-empty interior

Submodular and base polyhedra - Properties

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- Many facets (up to 2^p), many extreme points (up to p!)

Submodular and base polyhedra - Properties

- Submodular polyhedron: $P(F) = \{s \in \mathbb{R}^p, \ \forall A \subset V, \ s(A) \leqslant F(A)\}$
- Base polyhedron: $B(F) = P(F) \cap \{s(V) = F(V)\}$
- Many facets (up to 2^p), many extreme points (up to p!)
- \bullet Fundamental property (Edmonds, 1970): If F is submodular, maximizing linear functions may be done by a "greedy algorithm"
 - Let $w \in \mathbb{R}^p_+$ such that $w_{j_1} \geqslant \cdots \geqslant w_{j_p}$
 - Let $s_{j_k} = F(\{j_1, \dots, j_k\}) F(\{j_1, \dots, j_{k-1}\})$ for $k \in \{1, \dots, p\}$
 - Then $f(w) = \max_{s \in P(F)} w^{\top} s = \max_{s \in B(F)} w^{\top} s$
 - Both problems attained at s defined above
- Simple proof by convex duality

Submodular functions Links with convexity

• Theorem (Lovász, 1982): If F is submodular, then

$$\min_{A \subset V} F(A) = \min_{w \in \{0,1\}^p} f(w) = \min_{w \in [0,1]^p} f(w)$$

- Consequence: Submodular function minimization may be done in polynomial time (through ellipsoid algorithm)
- Representation of f(w) as a support function (Edmonds, 1970):

$$f(w) = \max_{s \in B(F)} s^{\top} w$$

- Maximizer s may be found efficiently through the greedy algorithm

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Submodular function minimization Dual problem

- Let $F: 2^V \to \mathbb{R}$ be a submodular function (such that $F(\varnothing) = 0$)
- Convex duality (Edmonds, 1970):

$$\min_{A \subset V} F(A) = \min_{w \in [0,1]^p} f(w)
= \min_{w \in [0,1]^p} \max_{s \in B(F)} w^{\top} s
= \max_{s \in B(F)} \min_{w \in [0,1]^p} w^{\top} s = \max_{s \in B(F)} s_{-}(V)$$

Exact submodular function minimization Combinatorial algorithms

- Algorithms based on $\min_{A \subset V} F(A) = \max_{s \in B(F)} s_{-}(V)$
- ullet Output the subset A and a base $s \in B(F)$ as a certificate of optimality
- Best algorithms have polynomial complexity (Schrijver, 2000; Iwata et al., 2001; Orlin, 2009) (typically $O(p^6)$ or more)
- Update a sequence of convex combination of vertices of B(F) obtained from the greedy algorithm using a specific order:
 - Based only on function evaluations
- Recent algorithms using efficient reformulations in terms of generalized graph cuts (Jegelka et al., 2011)

Approximate submodular function minimization

- For most machine learning applications, no need to obtain exact minimum
 - For convex optimization, see, e.g., Bottou and Bousquet (2008)

$$\min_{A \subset V} F(A) = \min_{w \in \{0,1\}^p} f(w) = \min_{w \in [0,1]^p} f(w)$$

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- ullet Important properties of f for convex optimization
 - Polyhedral function
 - Representation as maximum of linear functions

$$f(w) = \max_{s \in B(F)} w^{\top} s$$

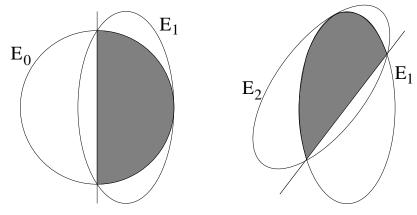
• Stability vs. speed vs. generality vs. ease of implementation

Projected subgradient descent (Shor et al., 1985)

- \bullet Subgradient of $f(w) = \max_{s \in B(F)} s^\top w$ through the greedy algorithm
- Using projected subgradient descent to minimize f on $[0,1]^p$
 - Iteration: $w_t = \prod_{[0,1]^p} \left(w_{t-1} \frac{C}{\sqrt{t}} s_t \right)$ where $s_t \in \partial f(w_{t-1})$
 - Convergence rate: $f(w_t) \min_{w \in [0,1]^p} f(w) \leq \frac{\sqrt{p}}{\sqrt{t}}$ with primal/dual guarantees (Nesterov, 2003)
- Fast iterations but slow convergence
 - need $O(p/\varepsilon^2)$ iterations to reach precision ε
 - need $O(p^2/\varepsilon^2)$ function evaluations to reach precision ε

Ellipsoid method (Nemirovski and Yudin, 1983)

 Build a sequence of minimum volume ellipsoids that enclose the set of solutions



- ullet Cost of a single iteration: p function evaluations and $O(p^3)$ operations
- Number of iterations: $2p^2 \left(\max_{A \subset V} F(A) \min_{A \subset V} F(A) \right) \log \frac{1}{\varepsilon}$.
 - $O(p^5)$ operations and $O(p^3)$ function evaluations
- Slow in practice (the bound is "tight")

Analytic center cutting planes (Goffin and Vial, 1993)

Center of gravity method

- improves the convergence rate of ellipsoid method
- cannot be computed easily
- Analytic center of a polytope defined by $a_i^\top w \leqslant b_i$, $i \in I$

$$\min_{w \in \mathbb{R}^p} - \sum_{i \in I} \log(b_i - a_i^\top w)$$

- Analytic center cutting planes (ACCPM)
 - Each iteration has complexity $O(p^2|I|+|I|^3)$ using Newton's method
 - No linear convergence rate
 - Good performance in practice

Simplex method for submodular minimization

- Mentioned by Girlich and Pisaruk (1997); McCormick (2005)
- Formulation as linear program: $s \in B(F) \Leftrightarrow s = S^{\top} \eta$, $S \in \mathbb{R}^{d \times p}$

$$\begin{split} \max_{s \in B(F)} s_{-}(V) &= \max_{\eta \geqslant 0, \ \eta^{\top} 1_d = 1} \sum_{i=1}^p \min\{(S^{\top} \eta)_i, 0\} \\ &= \max_{\eta \geqslant 0, \ \alpha \geqslant 0, \ \beta \geqslant 0} -\beta^{\top} 1_p \text{ such that } S^{\top} \eta - \alpha + \beta = 0, \ \eta^{\top} 1_d = 1. \end{split}$$

- ullet Column generation for simplex methods: only access the rows of S by maximizing linear functions
 - no complexity bound, may get global optimum if enough iterations

Separable optimization on base polyhedron

• Optimization of convex functions of the form $\Psi(w) + f(w)$ with f Lovász extension of F, and $\Psi(w) = \sum_{k \in V} \psi_k(w_k)$

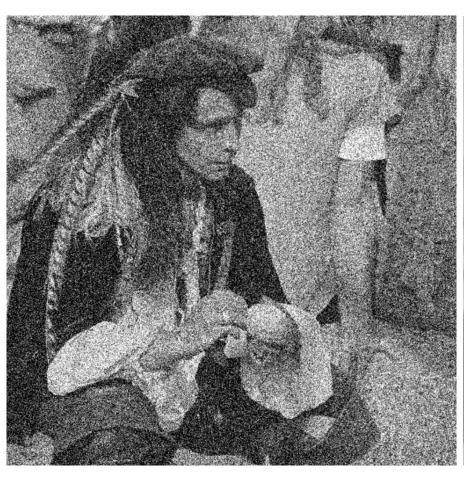
Structured sparsity

- Total variation denoising isotonic regression
- Regularized risk minimization penalized by the Lovász extension

Total variation denoising (Chambolle, 2005)

•
$$F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j)$$
 \Rightarrow $f(w) = \sum_{k, j \in V} d(k, j)(w_k - w_j)_+$

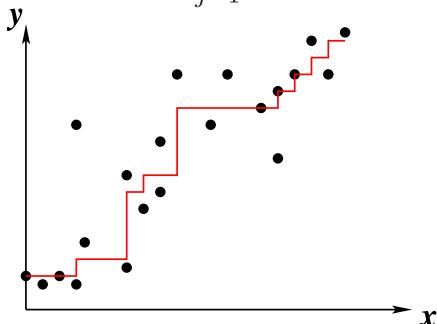
• d symmetric $\Rightarrow f = \text{total variation}$





Isotonic regression

- Given real numbers x_i , $i = 1, \ldots, p$
 - Find $y \in \mathbb{R}^p$ that minimizes $\frac{1}{2} \sum_{j=1}^p (x_i y_i)^2$ such that $\forall i, y_i \leqslant y_{i+1}$



- For a directed chain, f(y) = 0 if and only if $\forall i, y_i \leq y_{i+1}$
- Minimize $\frac{1}{2} \sum_{j=1}^{p} (x_i y_i)^2 + \lambda f(y)$ for λ large

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Separable optimization on base polyhedron

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Structured sparsity

- Total variation denoising isotonic regression
- Regularized risk minimization penalized by the Lovász extension
- Proximal methods (see second part)
 - Minimize $\Psi(w)+f(w)$ for smooth Ψ as soon as the following "proximal" problem may be obtained efficiently

$$\min_{w \in \mathbb{R}^p} \frac{1}{2} ||w - z||_2^2 + f(w) = \min_{w \in \mathbb{R}^p} \sum_{k=1}^p \frac{1}{2} (w_k - z_k)^2 + f(w)$$

• Submodular function minimization

Separable optimization on base polyhedron Convex duality

- Let $\psi_k : \mathbb{R} \to \mathbb{R}$, $k \in \{1, \dots, p\}$ be p functions. Assume
 - Each ψ_k is strictly convex
 - $-\sup_{\alpha\in\mathbb{R}}\psi_j'(\alpha)=+\infty \text{ and }\inf_{\alpha\in\mathbb{R}}\psi_j'(\alpha)=-\infty$
 - Denote ψ_1^*,\ldots,ψ_p^* their Fenchel-conjugates (then with full domain)

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$$\min_{w \in \mathbb{R}^{p}} f(w) + \sum_{j=1}^{p} \psi_{i}(w_{j}) = \min_{w \in \mathbb{R}^{p}} \max_{s \in B(F)} w^{\top} s + \sum_{j=1}^{p} \psi_{j}(w_{j})$$

$$= \max_{s \in B(F)} \min_{w \in \mathbb{R}^{p}} w^{\top} s + \sum_{j=1}^{p} \psi_{j}(w_{j})$$

$$= \max_{s \in B(F)} - \sum_{j=1}^{p} \psi_{j}^{*}(-s_{j})$$

Separable optimization on base polyhedron Equivalence with submodular function minimization

- For $\alpha \in \mathbb{R}$, let $A^{\alpha} \subset V$ be a minimizer of $A \mapsto F(A) + \sum_{j \in A} \psi'_j(\alpha)$
- Let w^* be the unique minimizer of $w \mapsto f(w) + \sum_{j=1}^p \psi_j(w_j)$
- Proposition (Chambolle and Darbon, 2009):
 - Given A^{α} for all $\alpha \in \mathbb{R}$, then $\forall j, \ w_j^* = \sup(\{\alpha \in \mathbb{R}, \ j \in A^{\alpha}\})$
 - Given w^* , then $A\mapsto F(A)+\sum_{j\in A}\psi_j'(\alpha)$ has minimal minimizer $\{w^*>\alpha\}$ and maximal minimizer $\{w^*\geqslant\alpha\}$
- Separable optimization equivalent to a sequence of submodular function minimizations
 - NB: extension of known results from parametric max-flow

Equivalence with submodular function minimization Proof sketch (Bach, 2011b)

• Duality gap for $\min_{w \in \mathbb{R}^p} f(w) + \sum_{j=1}^p \psi_i(w_j) = \max_{s \in B(F)} - \sum_{j=1}^p \psi_j^*(-s_j)$

$$f(w) + \sum_{j=1}^{p} \psi_{i}(w_{j}) - \sum_{j=1}^{p} \psi_{j}^{*}(-s_{j})$$

$$= f(w) - w^{\top}s + \sum_{j=1}^{p} \left\{ \psi_{j}(w_{j}) + \psi_{j}^{*}(-s_{j}) + w_{j}s_{j} \right\}$$

$$= \int_{-\infty}^{+\infty} \left\{ (F + \psi'(\alpha))(\{w \geqslant \alpha\}) - (s + \psi'(\alpha))_{-}(V) \right\} d\alpha$$

 Duality gap for convex problems = sums of duality gaps for combinatorial problems

Separable optimization on base polyhedron Quadratic case

- Let F be a submodular function and $w \in \mathbb{R}^p$ the unique minimizer of $w \mapsto f(w) + \frac{1}{2}||w||_2^2$. Then:
- (a) s=-w is the point in B(F) with minimum ℓ_2 -norm
- (b) For all $\lambda \in \mathbb{R}$, the maximal minimizer of $A \mapsto F(A) + \lambda |A|$ is $\{w \geqslant -\lambda\}$ and the minimal minimizer of F is $\{w > -\lambda\}$

Consequences

- Threshold at 0 the minimum norm point in B(F) to minimize F (Fujishige and Isotani, 2011)
- Minimizing submodular functions with cardinality constraints (Nagano et al., 2011)

From convex to combinatorial optimization

- Solving $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$ to solve $\min_{A \subset V} F(A)$
 - Thresholding solutions w at zero if $\forall k \in V, \psi'_k(0) = 0$
 - For quadratic functions $\psi_k(w_k) = \frac{1}{2}w_k^2$, equivalent to projecting 0 on B(F) (Fujishige, 2005)

From convex to combinatorial optimization and vice-versa...

- Solving $\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$ to solve $\min_{A \subset V} F(A)$
 - Thresholding solutions w at zero if $\forall k \in V, \psi_k'(0) = 0$
 - For quadratic functions $\psi_k(w_k)=\frac{1}{2}w_k^2$, equivalent to projecting 0 on B(F) (Fujishige, 2005)
- Solving $\min_{A\subset V}F(A)-t(A)$ to solve $\min_{w\in\mathbb{R}^p}\sum_{k\in V}\psi_k(w_k)+f(w)$
 - General decomposition strategy (Groenevelt, 1991)
 - Efficient only when submodular minimization is efficient

Solving
$$\min_{A\subset V}F(A)-t(A)$$
 to solve $\min_{w\in\mathbb{R}^p}\sum_{k\in V}\psi_k(w_k)+f(w)$

- General recursive divide-and-conquer algorithm (Groenevelt, 1991)
- NB: Dual version of Fujishige (2005)
 - 1. Compute minimizer $t \in \mathbb{R}^p$ of $\sum_{j \in V} \psi_j^*(-t_j)$ s.t. t(V) = F(V)
 - 2. Compute minimizer A of F(A) t(A)
 - 3. If A = V, then t is optimal. Exit.
 - 4. Compute a minimizer s_A of $\sum_{j\in A} \psi_j^*(-s_j)$ over $s\in B(F_A)$ where $F_A:2^A\to\mathbb{R}$ is the restriction of F to A, i.e., $F_A(B)=F(A)$
 - 5. Compute a minimizer $s_{V\setminus A}$ of $\sum_{j\in V\setminus A}\psi_j^*(-s_j)$ over $s\in B(F^A)$ where $F^A(B)=F(A\cup B)-F(A)$, for $B\subset V\setminus A$
 - 6. Concatenate s_A and $s_{V \setminus A}$. Exit.

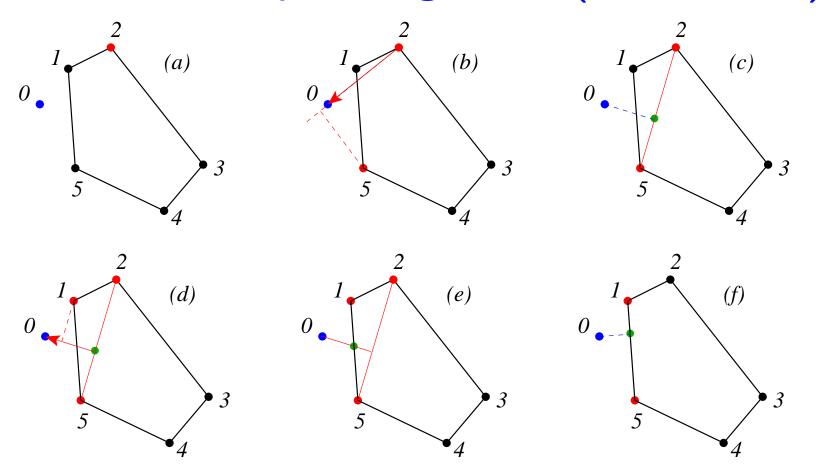
Solving
$$\min_{w \in \mathbb{R}^p} \sum_{k \in V} \psi_k(w_k) + f(w)$$
 to solve $\min_{A \subset V} F(A)$

- Dual problem: $\max_{s \in B(F)} \sum_{j=1}^{p} \psi_j^*(-s_j)$
- Constrained optimization when linear functions can be maximized
 - Frank-Wolfe algorithms
- Two main types for convex functions

• Goal:
$$\min_{w \in \mathbb{R}^p} \frac{1}{2} \|w\|_2^2 + f(w) = \max_{s \in B(F)} -\frac{1}{2} \|s\|_2^2$$

- ullet Can only maximize linear functions on B(F)
- Two types of "Frank-wolfe" algorithms
- 1. Active set algorithm (⇔ min-norm-point)
 - Sequence of maximizations of linear functions over B(F) + overheads (affine projections)
 - Finite convergence, but no complexity bounds

Minimum-norm-point algorithm (Wolfe, 1976)



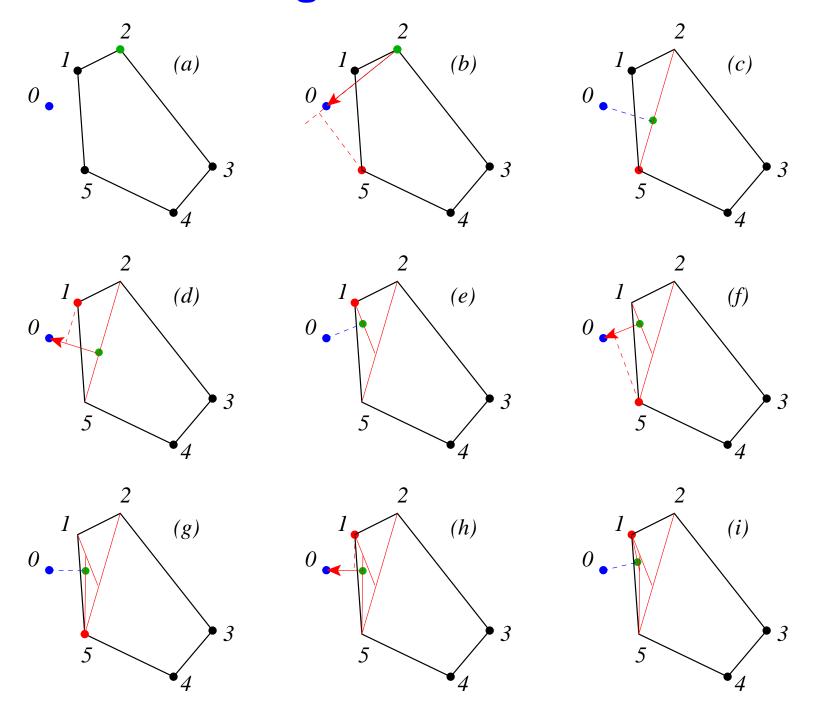
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• 2. Conditional gradient

- Sequence of maximizations of linear functions over B(F)
- Approximate optimality bound

Conditional gradient with line search



• **Proposition**: t steps of conditional gradient (with line search) outputs $s_t \in B(F)$ and $w_t = -s_t$, such that

$$f(w_t) + \frac{1}{2} ||w_t||_2^2 - \text{OPT} \leqslant f(w_t) + \frac{1}{2} ||w_t||_2^2 + \frac{1}{2} ||s_t||_2^2 \leqslant \frac{2D^2}{t}$$

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- Improved primal candidate through isotonic regression
 - -f(w) is linear on any set of w with fixed ordering
 - May be optimized using isotonic regression ("pool-adjacent-violator") in O(n) (see, e.g., Best and Chakravarti, 1990)
 - Given $w_t = -s_t$, keep the ordering and reoptimize

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 - Given $w_t = -s_t$, keep the ordering and reoptimize
- Better bound for submodular function minimization?

From quadratic optimization on B(F) to submodular function minimization

- **Proposition**: If w is ε -optimal for $\min_{w \in \mathbb{R}^p} \frac{1}{2} ||w||_2^2 + f(w)$, then at least a levet set A of w is $\left(\frac{\sqrt{\varepsilon p}}{2}\right)$ -optimal for submodular function minimization
- If $\varepsilon=\frac{2D^2}{t}$, $\frac{\sqrt{\varepsilon p}}{2}=\frac{Dp^{1/2}}{\sqrt{2t}}\Rightarrow$ no provable gains, but:
 - Bound on the iterates A_t (with additional assumptions)
 - Possible thresolding for acceleration

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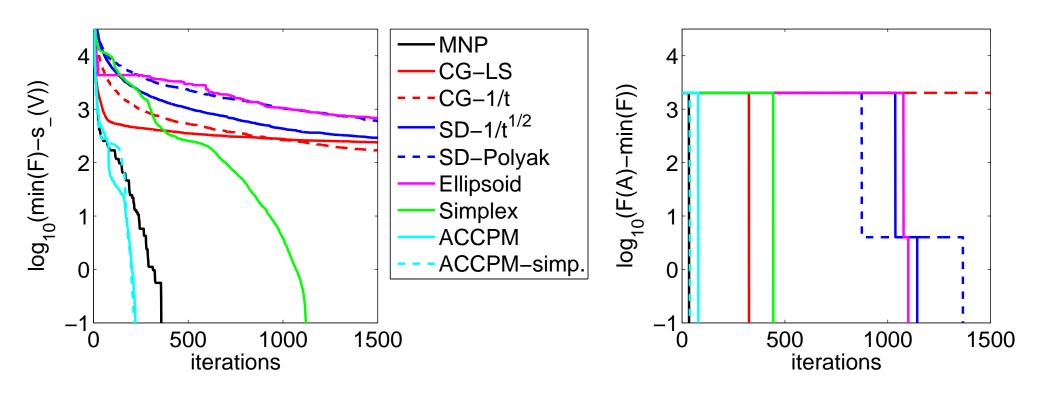
• If
$$\varepsilon=\frac{2D^2}{t}$$
, $\frac{\sqrt{\varepsilon p}}{2}=\frac{Dp^{1/2}}{\sqrt{2t}}\Rightarrow$ no provable gains, but:

- Bound on the iterates A_t (with additional assumptions)
- Possible thresolding for acceleration
- Lower complexity bound for SFM
 - Conjecture: no algorithm that is based only on a sequence of greedy algorithms obtained from linear combinations of bases can improve on the subgradient bound (after p/2 iterations).

Simulations on standard benchmark "DIMACS Genrmf-wide", p = 430

Submodular function minimization

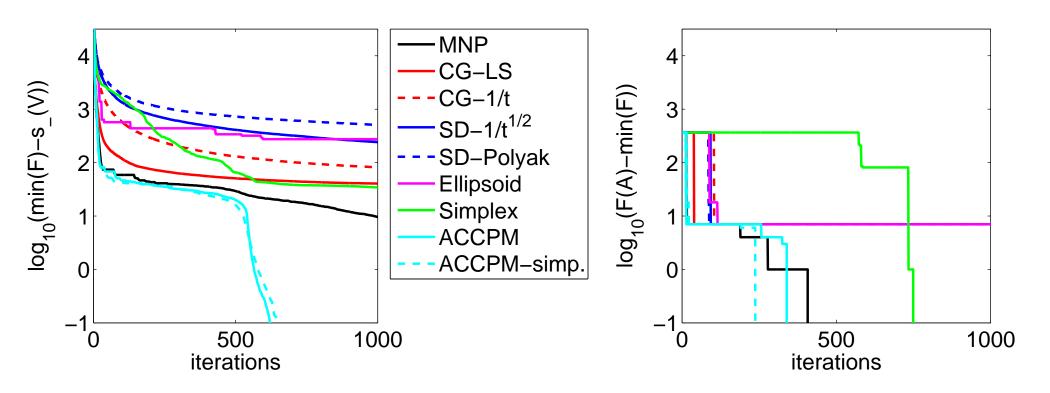
- (Left) dual suboptimality
- (Right) primal suboptimality



Simulations on standard benchmark "DIMACS Genrmf-long", p = 575

Submodular function minimization

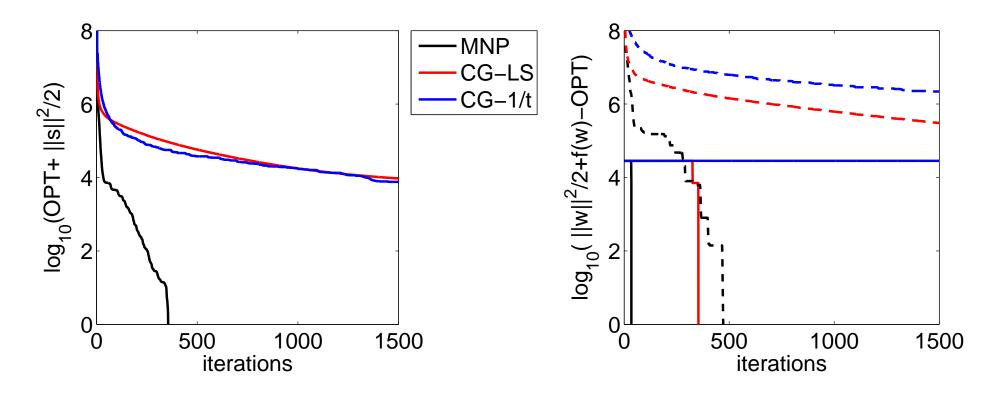
- (Left) dual suboptimality
- (Right) primal suboptimality



Simulations on standard benchmark

Separable quadratic optimization

- (Left) dual suboptimality
- (Right) primal suboptimality
 (in dashed, before the pool-adjacent-violator correction)



Outline

1. Submodular functions

- Review and examples of submodular functions
- Links with convexity through Lovász extension

2. Submodular minimization

- Non-smooth convex optimization
- Parallel algorithm for special case

3. Structured sparsity-inducing norms

- Relaxation of the penalization of supports by submodular functions
- Extensions (symmetric, ℓ_q -relaxation)

From submodular minimization to proximal problems

- Summary: several optimization problems
 - Discrete problem: $\min_{A \subset V} F(A) = \min_{w \in \{0,1\}^p} f(w)$
 - Continuous problem: $\min_{w \in [0,1]^p} f(w)$
 - Proximal problem (P): $\min_{w \in \mathbb{R}^p} \frac{1}{2} ||w||_2^2 + f(w)$
- ullet Solving (P) is equivalent to minimizing $F(A) + \lambda |A|$ for all λ

$$-\arg\min_{A\subseteq V} F(A) + \lambda |A| = \{k, w_k \geqslant -\lambda\}$$

- Much simpler problem but no gains in terms of (provable) complexity
 - See Bach (2011a)

Decomposable functions

 \bullet F may often be decomposed as the sum of r "simple" functions:

$$F(A) = \sum_{j=1}^{r} F_j(A)$$

- Each F_i may be minimized efficiently
- Example: 2D grid = vertical chains + horizontal chains
- Komodakis et al. (2011); Kolmogorov (2012); Stobbe and Krause (2010); Savchynskyy et al. (2011)
 - Dual decomposition approach but slow non-smooth problem

Decomposable functions and proximal problems (Jegelka, Bach, and Sra, 2013)

Dual problem

$$\min_{w \in \mathbb{R}^p} f_1(w) + f_2(w) + \frac{1}{2} \|w\|_2^2$$

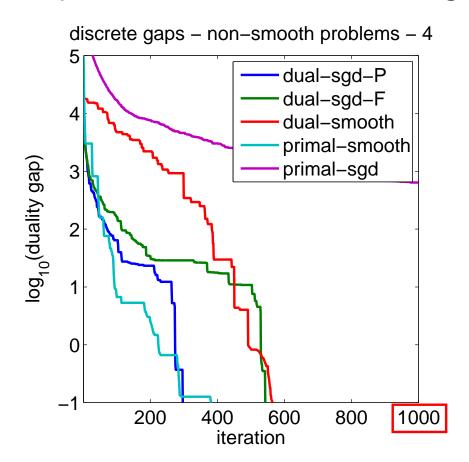
$$= \min_{w \in \mathbb{R}^p} \max_{s_1 \in B(F_1)} s_1^\top w + \max_{s_2 \in B(F_2)} s_2^\top w + \frac{1}{2} \|w\|_2^2$$

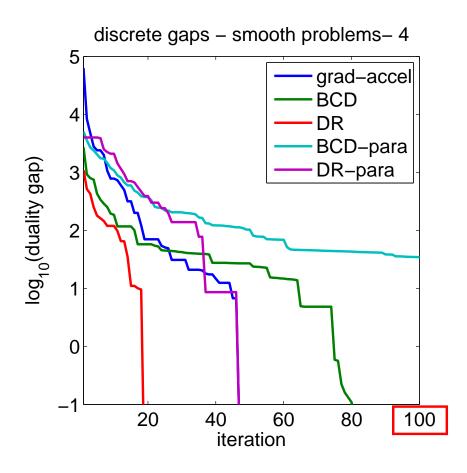
$$= \max_{s_1 \in B(F_1), \ s_2 \in B(F_2)} -\frac{1}{2} \|s_1 + s_2\|^2$$

- Finding the closest point between two polytopes
 - Several alternatives: Block coordinate ascent, Douglas Rachford splitting (Bauschke et al., 2004)
 - (a) no parameters, (b) parallelizable

Experiments

 \bullet Graph cuts on a 500×500 image

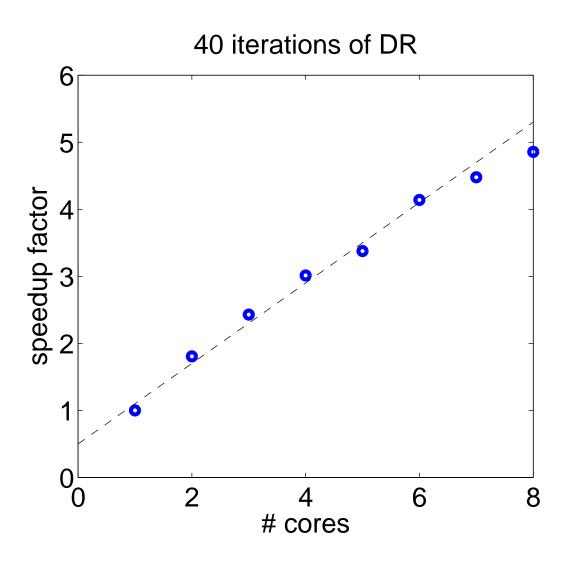




- Matlab/C implementation 10 times slower than C-code for graph cut
 - Easy to code and parallelizable

Parallelization

• Multiple cores



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Structured sparsity through submodular functions References and Links

References on submodular functions

- Submodular Functions and Optimization (Fujishige, 2005)
- Tutorial paper based on convex optimization (Bach, 2011b)
 www.di.ens.fr/~fbach/submodular_fot.pdf

Structured sparsity through convex optimization

- Algorithms (Bach, Jenatton, Mairal, and Obozinski, 2011) www.di.ens.fr/~fbach/bach_jenatton_mairal_obozinski_FOT.pdf
- Theory/applications (Bach, Jenatton, Mairal, and Obozinski, 2012)
 www.di.ens.fr/~fbach/stat_science_structured_sparsity.pdf
- Matlab/R/Python codes: http://www.di.ens.fr/willow/SPAMS/
- Slides: www.di.ens.fr/~fbach/fbach_cargese_2013.pdf

Sparsity in supervised machine learning

- Observed data $(x_i, y_i) \in \mathbb{R}^p \times \mathbb{R}$, $i = 1, \ldots, n$
 - Response vector $y = (y_1, \dots, y_n)^{\top} \in \mathbb{R}^n$
 - Design matrix $X = (x_1, \dots, x_n)^{\top} \in \mathbb{R}^{n \times p}$
- Regularized empirical risk minimization:

$$\min_{w \in \mathbb{R}^p} \frac{1}{n} \sum_{i=1}^n \ell(y_i, w^\top x_i) + \lambda \Omega(w) = \left[\min_{w \in \mathbb{R}^p} L(y, Xw) + \lambda \Omega(w) \right]$$

- Norm Ω to promote sparsity
 - square loss + ℓ_1 -norm \Rightarrow basis pursuit in signal processing (Chen et al., 2001), Lasso in statistics/machine learning (Tibshirani, 1996)
 - Proxy for interpretability
 - Allow high-dimensional inference: $\log p = O(n)$

Sparsity in unsupervised machine learning

• Multiple responses/signals $y = (y^1, \dots, y^k) \in \mathbb{R}^{n \times k}$

$$\min_{w^1, \dots, w^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(y^j, Xw^j) + \lambda \Omega(w^j) \right\}$$

Sparsity in unsupervised machine learning

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- Only responses are observed ⇒ Dictionary learning
 - Learn $X=(x^1,\ldots,x^p)\in\mathbb{R}^{n\times p}$ such that $\forall j,\ \|x^j\|_2\leqslant 1$

$$\min_{X=(x^1,...,x^p)} \min_{w^1,...,w^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(y^j, Xw^j) + \lambda \Omega(w^j) \right\}$$

- Olshausen and Field (1997); Elad and Aharon (2006); Mairal et al.
 (2009a)
- sparse PCA: replace $||x^j||_2 \leqslant 1$ by $\Theta(x^j) \leqslant 1$

Sparsity in signal processing

• Multiple responses/signals $x = (x^1, \dots, x^k) \in \mathbb{R}^{n \times k}$

$$\min_{\alpha^1, \dots, \alpha^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(x^j, D\alpha^j) + \lambda \Omega(\alpha^j) \right\}$$

- Only responses are observed ⇒ Dictionary learning
 - Learn $D=(d^1,\ldots,d^p)\in\mathbb{R}^{n\times p}$ such that $\forall j,\ \|d^j\|_2\leqslant 1$

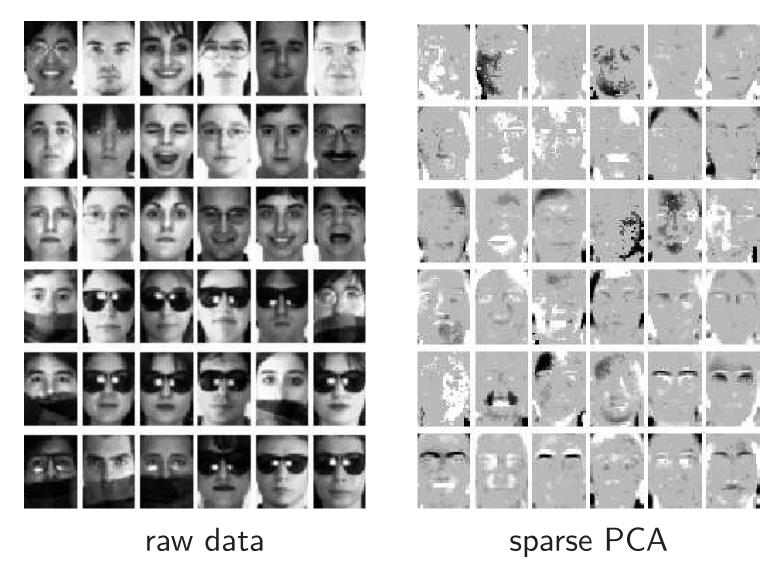
$$\min_{D=(d^1,\dots,d^p)} \min_{\alpha^1,\dots,\alpha^k \in \mathbb{R}^p} \sum_{j=1}^k \left\{ L(x^j, D\alpha^j) + \lambda \Omega(\alpha^j) \right\}$$

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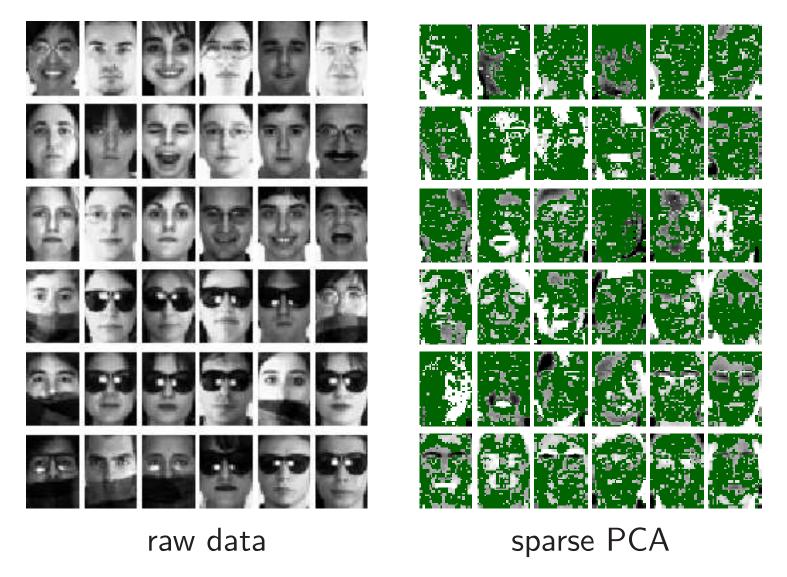
Why structured sparsity?

Interpretability

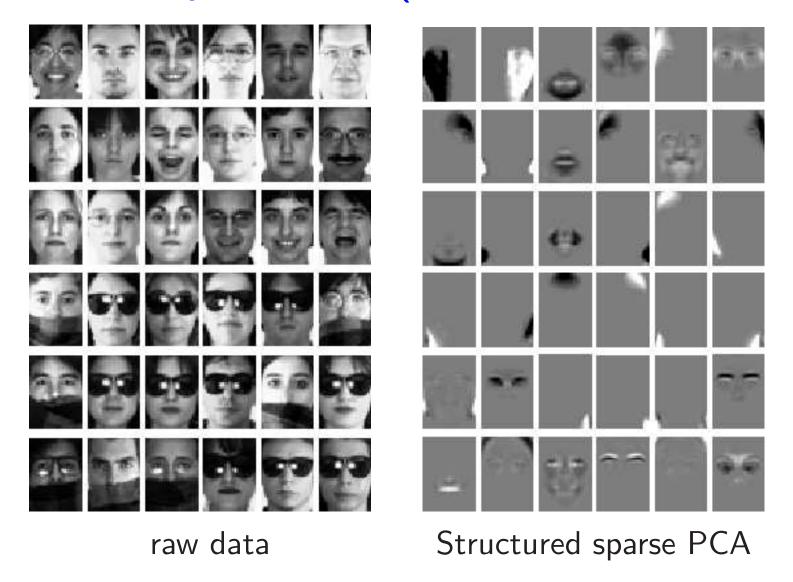
- Structured dictionary elements (Jenatton et al., 2009b)
- Dictionary elements "organized" in a tree or a grid (Kavukcuoglu et al., 2009; Jenatton et al., 2010; Mairal et al., 2010)



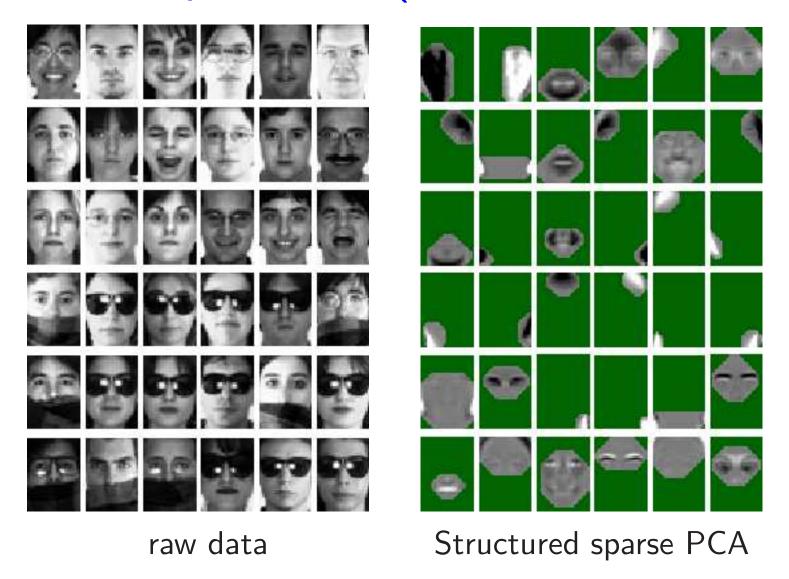
ullet Unstructed sparse PCA \Rightarrow many zeros do not lead to better interpretability



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ullet Enforce selection of convex nonzero patterns \Rightarrow robustness to occlusion in face identification



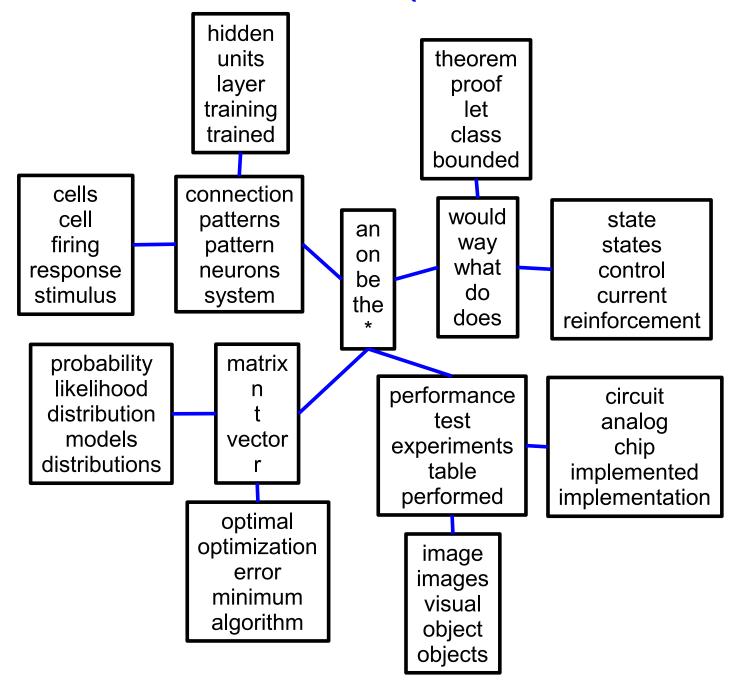
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Modelling of text corpora (Jenatton et al., 2010)



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Stability and identifiability

Prediction or estimation performance

When prior knowledge matches data (Haupt and Nowak, 2006;
 Baraniuk et al., 2008; Jenatton et al., 2009a; Huang et al., 2009)

Numerical efficiency

- Non-linear variable selection with 2^p subsets (Bach, 2008)

Classical approaches to structured sparsity

Many application domains

- Computer vision (Cevher et al., 2008; Mairal et al., 2009b)
- Neuro-imaging (Gramfort and Kowalski, 2009; Jenatton et al., 2011)
- Bio-informatics (Rapaport et al., 2008; Kim and Xing, 2010)

Non-convex approaches

Haupt and Nowak (2006); Baraniuk et al. (2008); Huang et al. (2009)

Convex approaches

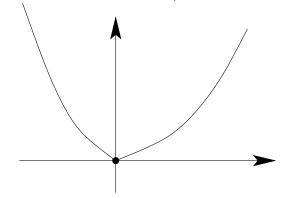
Design of sparsity-inducing norms

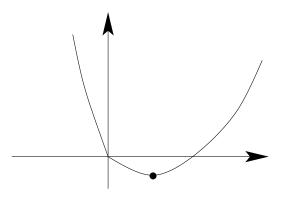
Why ℓ_1 -norms lead to sparsity?

Example 1: quadratic problem in 1D, i.e., $\left| \min_{x \in \mathbb{R}} \frac{1}{2} x^2 - xy + \lambda |x| \right|$

$$\min_{x \in \mathbb{R}} \frac{1}{2} x^2 - xy + \lambda |x|$$

- Piecewise quadratic function with a kink at zero
 - Derivative at $0+: g_+ = \lambda y$ and $0-: g_- = -\lambda y$





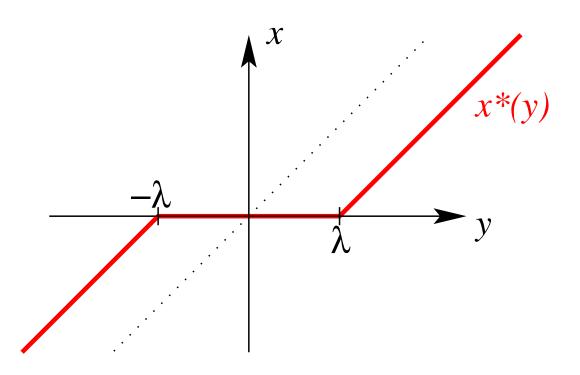
- -x=0 is the solution iff $g_{+}\geqslant 0$ and $g_{-}\leqslant 0$ (i.e., $|y|\leqslant \lambda$)
- $-x \geqslant 0$ is the solution iff $g_+ \leqslant 0$ (i.e., $y \geqslant \lambda$) $\Rightarrow x^* = y \lambda$
- $-x \leqslant 0$ is the solution iff $g_{-} \leqslant 0$ (i.e., $y \leqslant -\lambda$) $\Rightarrow x^* = y + \lambda$
- Solution $|x^* = \operatorname{sign}(y)(|y| \lambda)_+| = \operatorname{soft\ thresholding}$

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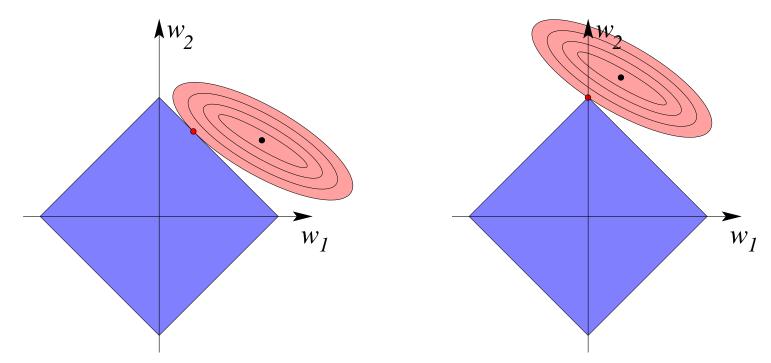
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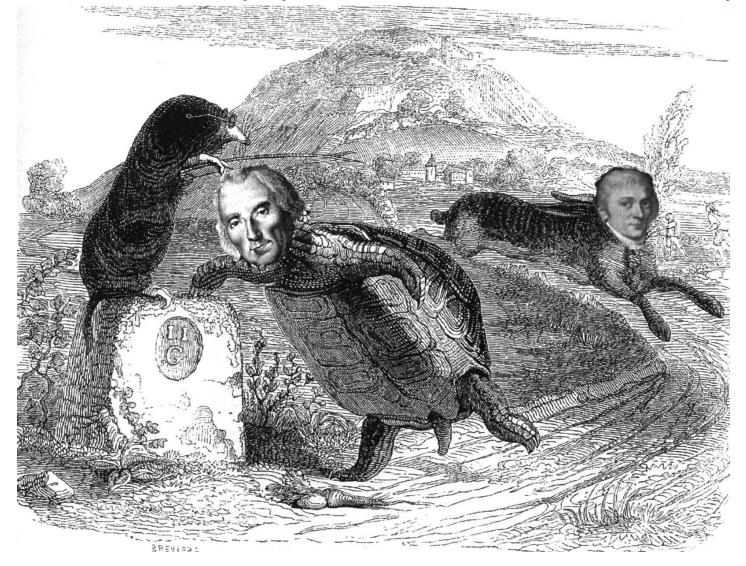
Why ℓ_1 -norms lead to sparsity?

- **Example 2**: minimize quadratic function Q(w) subject to $||w||_1 \leqslant T$.
 - coupled soft thresholding
- Geometric interpretation
 - NB: penalizing is "equivalent" to constraining



Non-smooth optimization!

Gaussian hare (ℓ_2) vs. Laplacian tortoise (ℓ_1)



- Smooth vs. non-smooth optimization
- See Bach, Jenatton, Mairal, and Obozinski (2011)

Sparsity-inducing norms

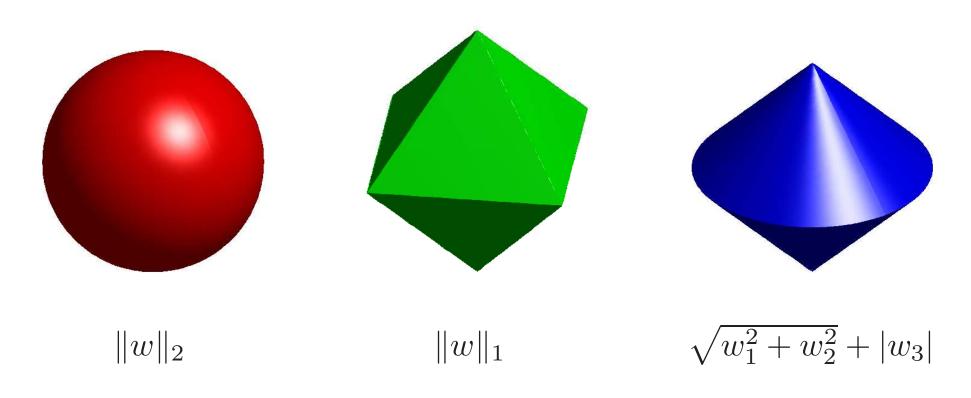
- Popular choice for Ω
 - The ℓ_1 - ℓ_2 norm,

$$\sum_{G \in \mathbf{H}} ||w_G||_2 = \sum_{G \in \mathbf{H}} \left(\sum_{j \in G} w_j^2\right)^{1/2}$$

- with \mathbf{H} a partition of $\{1,\ldots,p\}$
- The ℓ_1 - ℓ_2 norm sets to zero groups of non-overlapping variables (as opposed to single variables for the ℓ_1 -norm)
- For the square loss, group Lasso (Yuan and Lin, 2006)



Unit norm balls Geometric interpretation



Sparsity-inducing norms

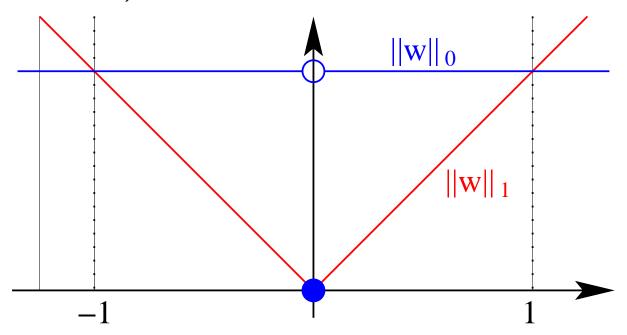
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- The ℓ_1 - ℓ_2 norm sets to zero groups of non-overlapping variables (as opposed to single variables for the ℓ_1 -norm)
- For the square loss, group Lasso (Yuan and Lin, 2006)
- What if the set of groups H is not a partition anymore?
- Is there any systematic way?

ℓ_1 -norm = convex envelope of cardinality of support

- Let $w \in \mathbb{R}^p$. Let $V = \{1, \dots, p\}$ and $\mathrm{Supp}(w) = \{j \in V, \ w_j \neq 0\}$
- Cardinality of support: $||w||_0 = \operatorname{Card}(\operatorname{Supp}(w))$
- Convex envelope = largest convex lower bound (see, e.g., Boyd and Vandenberghe, 2004)



• ℓ_1 -norm = convex envelope of ℓ_0 -quasi-norm on the ℓ_∞ -ball $[-1,1]^p$

Convex envelopes of general functions of the support (Bach, 2010)

- Let $F: 2^V \to \mathbb{R}$ be a **set-function**
 - Assume F is **non-decreasing** (i.e., $A \subset B \Rightarrow F(A) \leqslant F(B)$)
 - Explicit prior knowledge on supports (Haupt and Nowak, 2006;
 Baraniuk et al., 2008; Huang et al., 2009)
- Define $\Theta(w) = F(\operatorname{Supp}(w))$: How to get its convex envelope?
 - 1. Possible if F is also **submodular**
 - 2. Allows unified theory and algorithm
 - 3. Provides **new** regularizers

Submodular functions and structured sparsity

- ullet Let $F:2^V o \mathbb{R}$ be a non-decreasing submodular set-function
- Proposition: the convex envelope of $\Theta: w \mapsto F(\operatorname{Supp}(w))$ on the ℓ_{∞} -ball is $\Omega: w \mapsto f(|w|)$ where f is the Lovász extension of F

Proof - I

- Notation: $g: w \mapsto F(\operatorname{supp}(w))$ defined on $[-1,1]^p$
- Computation of the Fenchel dual

$$\begin{split} g^*(s) &= \max_{\|w\|_{\infty} \leqslant 1} w^\top s - g(w) \\ &= \max_{\delta \in \{0,1\}^p} \max_{\|w\|_{\infty} \leqslant 1} (\delta \circ w)^\top s - f(\delta) \text{ by definition of } g \\ &= \max_{\delta \in \{0,1\}^p} \delta^\top |s| - f(\delta) \text{ by maximizing out } w \\ &= \max_{\delta \in [0,1]^p} \delta^\top |s| - f(\delta) \text{ because } F - |s| \text{ is submodular} \end{split}$$

Proof - II

• Notation: $g: w \mapsto F(\operatorname{supp}(w))$ defined on $[-1,1]^p$

• Fenchel dual: $g^*(s) = \max_{\delta \in [0,1]^p} \delta^\top |s| - f(\delta)$

Proof - II

- Notation: $g: w \mapsto F(\operatorname{supp}(w))$ defined on $[-1,1]^p$
- Fenchel dual: $g^*(s) = \max_{\delta \in [0,1]^p} \delta^\top |s| f(\delta)$
- Computation of the Fenchel bi-dual, for all w such that $||w||_{\infty} \leq 1$:

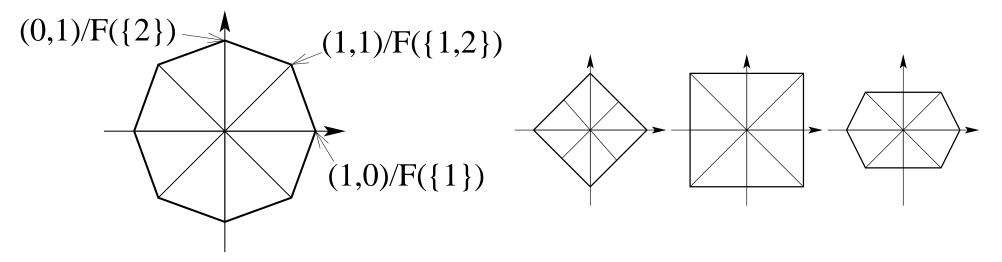
$$\begin{split} g^{**}(w) &= \max_{s \in \mathbb{R}^p} s^\top w - g^*(s) \\ &= \max_{s \in \mathbb{R}^p} \min_{\delta \in [0,1]^p} s^\top w - \delta^\top |s| + f(\delta) \\ &= \min_{\delta \in [0,1]^p} \max_{s \in \mathbb{R}^p} s^\top w - \delta^\top |s| + f(\delta) \text{ by strong duality} \\ &= \min_{\delta \in [0,1]^p, \delta \geqslant |w|} f(\delta) = f(|w|) \text{ because } F \text{ is nonincreasing} \end{split}$$

Submodular functions and structured sparsity

- ullet Let $F:2^V o \mathbb{R}$ be a non-decreasing submodular set-function
- Proposition: the convex envelope of $\Theta: w \mapsto F(\operatorname{Supp}(w))$ on the ℓ_{∞} -ball is $\Omega: w \mapsto f(|w|)$ where f is the Lovász extension of F

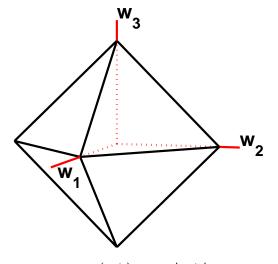
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- ullet Sparsity-inducing properties: Ω is a polyhedral norm



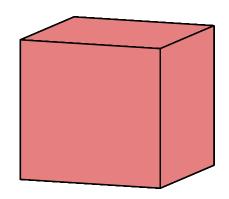
- A if stable if for all $B \supset A$, $B \neq A \Rightarrow F(B) > F(A)$
- With probability one, stable sets are the only allowed active sets

Polyhedral unit balls

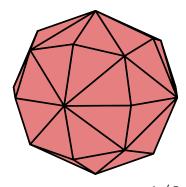


$$F(A) = |A|$$

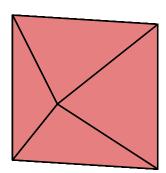
$$\Omega(w) = ||w||_1$$



 $F(A) = \min\{|A|, 1\}$ $\Omega(w) = ||w||_{\infty}$

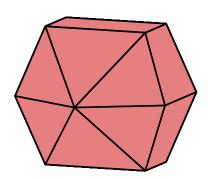


 $F(A) = |A|^{1/2}$ all possible extreme points



$$F(A) = 1_{\{A \cap \{1\} \neq \emptyset\}} + 1_{\{A \cap \{2,3\} \neq \emptyset\}}$$

$$\Omega(w) = |w_1| + ||w_{\{2,3\}}||_{\infty}$$



$$F(A) = 1_{\{A \cap \{1,2,3\} \neq \varnothing\}}$$

$$+1_{\{A \cap \{2,3\} \neq \varnothing\}} + 1_{\{A \cap \{3\} \neq \varnothing\}}$$

$$\Omega(w) = ||w||_{\infty} + ||w_{\{2,3\}}||_{\infty} + |w_{3}|$$

Submodular functions and structured sparsity Examples

- From $\Omega(w)$ to F(A): provides new insights into existing norms
 - Grouped norms with **overlapping** groups (Jenatton et al., 2009a)

$$\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty}$$

- $-\ell_1$ - ℓ_∞ norm \Rightarrow sparsity at the group level
- Some w_G 's are set to zero for some groups G

$$\left(\operatorname{Supp}(w)\right)^c = \bigcup_{G \in \mathbf{H}'} G \text{ for some } \mathbf{H}' \subseteq \mathbf{H}$$

Submodular functions and structured sparsity Examples

- From $\Omega(w)$ to F(A): provides new insights into existing norms
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$$\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty} \implies F(A) = \operatorname{Card}(\{G \in \mathbf{H}, G \cap A \neq \emptyset\})$$

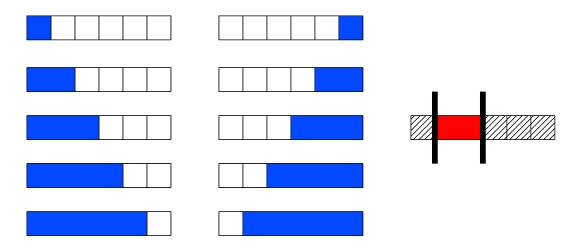
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Justification not only limited to allowed sparsity patterns

Selection of contiguous patterns in a sequence

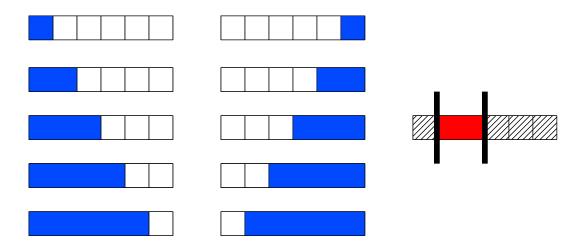
• Selection of contiguous patterns in a sequence



• H is the set of blue groups: any union of blue groups set to zero leads to the selection of a **contiguous pattern**

Selection of contiguous patterns in a sequence

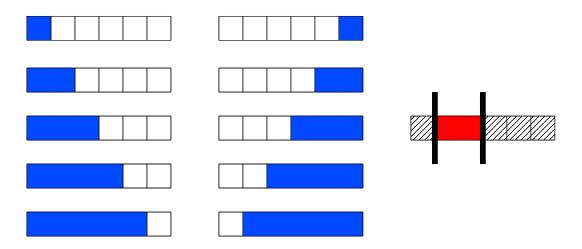
Selection of contiguous patterns in a sequence



- H is the set of blue groups: any union of blue groups set to zero leads to the selection of a **contiguous pattern**
- $\sum_{G \in \mathbf{H}} \|w_G\|_{\infty} \Rightarrow F(A) = p 2 + \operatorname{Range}(A) \text{ if } A \neq \emptyset$

Examples of set of groups H

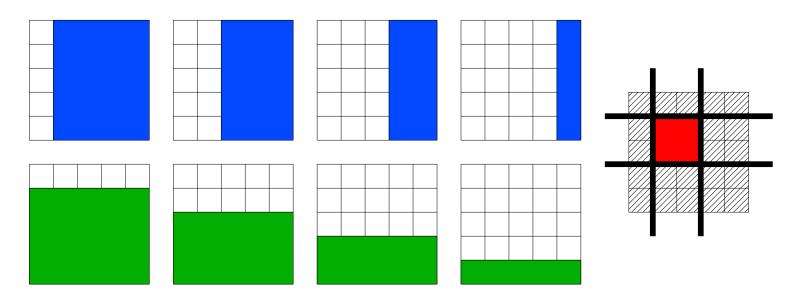
ullet Selection of contiguous patterns on a sequence, p=6



- H is the set of blue groups
- Any union of blue groups set to zero leads to the selection of a contiguous pattern

Examples of set of groups H

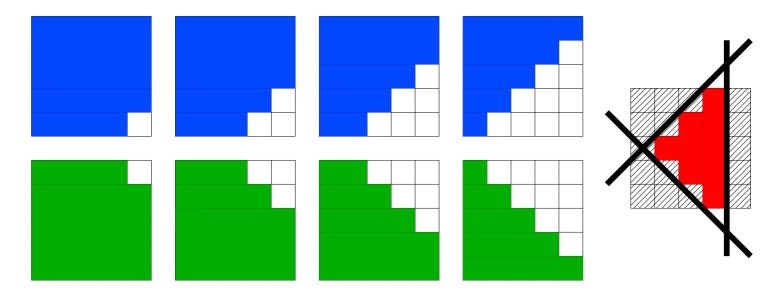
ullet Selection of rectangles on a 2-D grids, p=25



- H is the set of blue/green groups (with their not displayed complements)
- Any union of blue/green groups set to zero leads to the selection of a rectangle

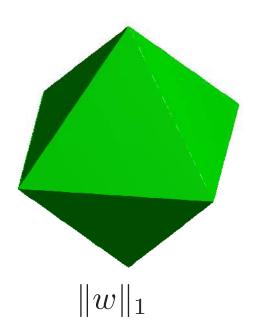
Examples of set of groups H

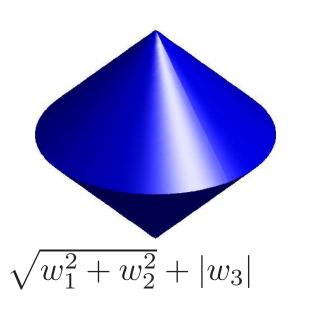
ullet Selection of diamond-shaped patterns on a 2-D grids, p=25.

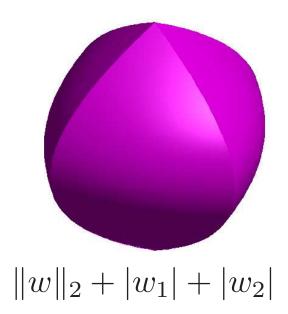


 It is possible to extend such settings to 3-D space, or more complex topologies

Unit norm balls Geometric interpretation







Application to background subtraction (Mairal, Jenatton, Obozinski, and Bach, 2010)

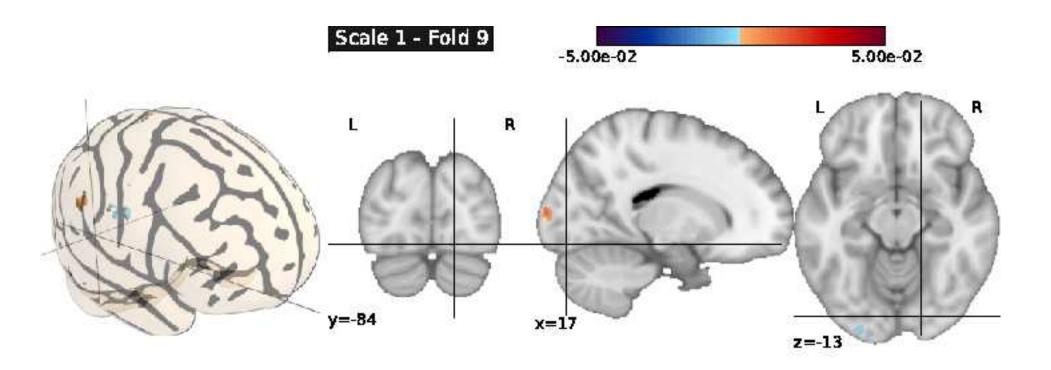
Input ℓ_1 -norm Structured norm

Application to background subtraction (Mairal, Jenatton, Obozinski, and Bach, 2010)

Background ℓ_1 -norm Structured norm

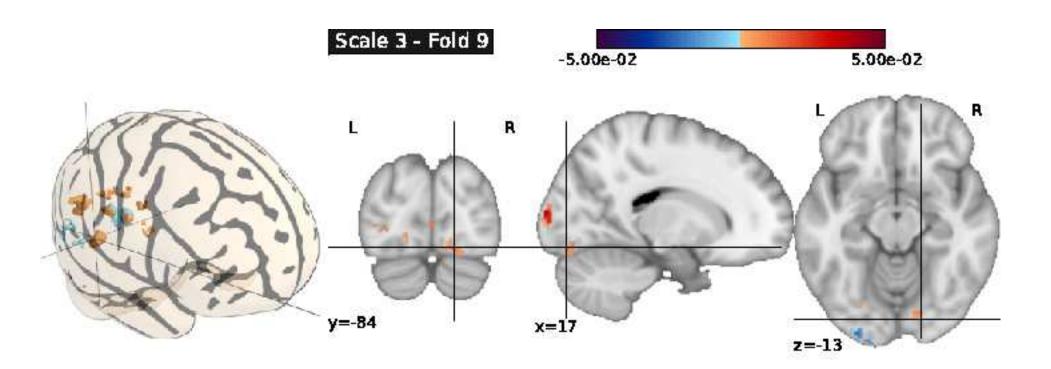
Application to neuro-imaging Structured sparsity for fMRI (Jenatton et al., 2011)

- "Brain reading": prediction of (seen) object size
- Multi-scale activity levels through hierarchical penalization



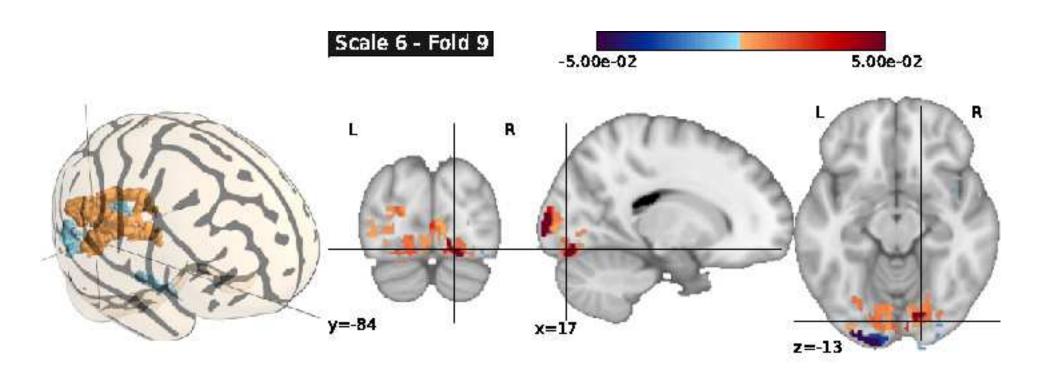
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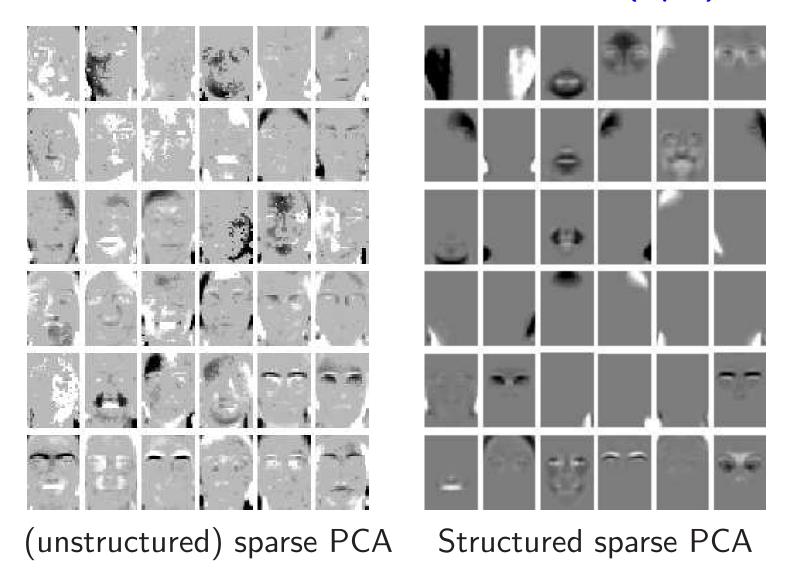


Sparse Structured PCA (Jenatton, Obozinski, and Bach, 2009b)

• Learning sparse and structured dictionary elements:

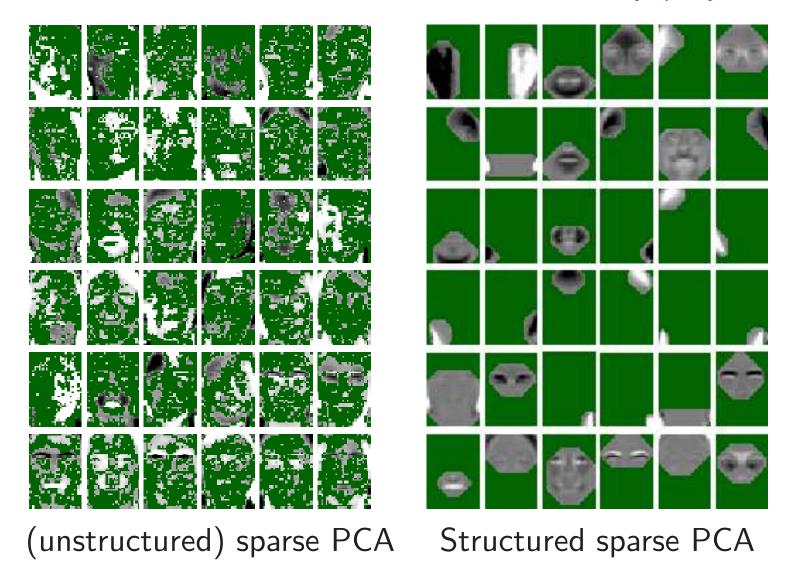
$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^{n} \|y^i - Xw^i\|_2^2 + \lambda \sum_{j=1}^{p} \Omega(x^j) \text{ s.t. } \forall i, \ \|w^i\|_2 \leq 1$$

Application to face databases (2/3)



ullet Enforce selection of convex nonzero patterns \Rightarrow robustness to occlusion

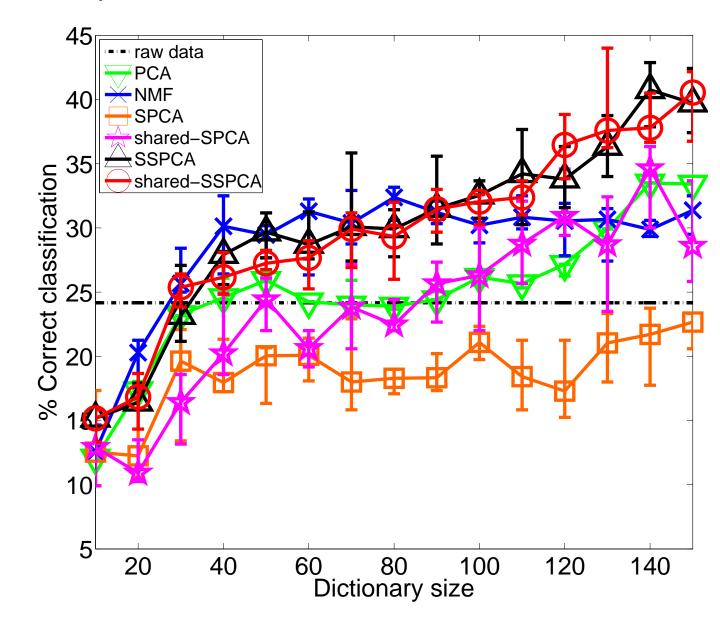
Application to face databases (2/3)



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Application to face databases (3/3)

• Quantitative performance evaluation on classification task



Dictionary learning vs. sparse structured PCA Exchange roles of X and w

• Sparse structured PCA (structured dictionary elements):

$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^n \|y^i - Xw^i\|_2^2 + \lambda \sum_{j=1}^k \Omega(x^j) \text{ s.t. } \forall i, \ \|w^i\|_2 \, \leq \, 1.$$

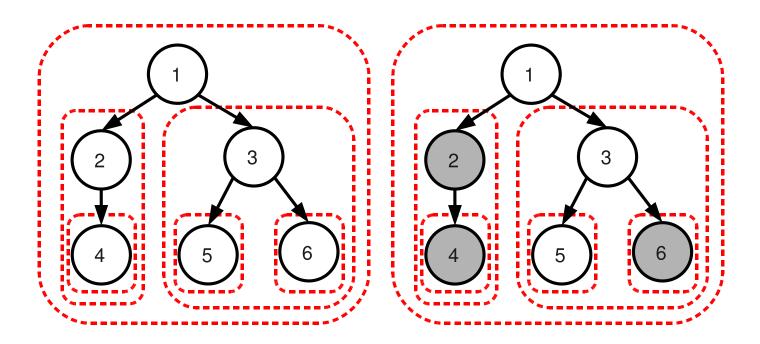
ullet Dictionary learning with **structured sparsity for codes** w:

$$\min_{W \in \mathbb{R}^{k \times n}, X \in \mathbb{R}^{p \times k}} \frac{1}{n} \sum_{i=1}^{n} \|y^i - Xw^i\|_2^2 + \lambda \Omega(w^i) \text{ s.t. } \forall j, \ \|x^j\|_2 \, \leq \, 1.$$

- Optimization: proximal methods
 - Requires solving many times $\min_{w \in \mathbb{R}^p} \frac{1}{2} \|y w\|_2^2 + \lambda \Omega(w)$
 - Modularity of implementation if proximal step is efficient (Jenatton et al., 2010; Mairal et al., 2010)

Hierarchical dictionary learning (Jenatton, Mairal, Obozinski, and Bach, 2010)

- Structure on codes w (not on dictionary X)
- Hierarchical penalization: $\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty}$ where groups G in \mathbf{H} are equal to set of descendants of some nodes in a tree

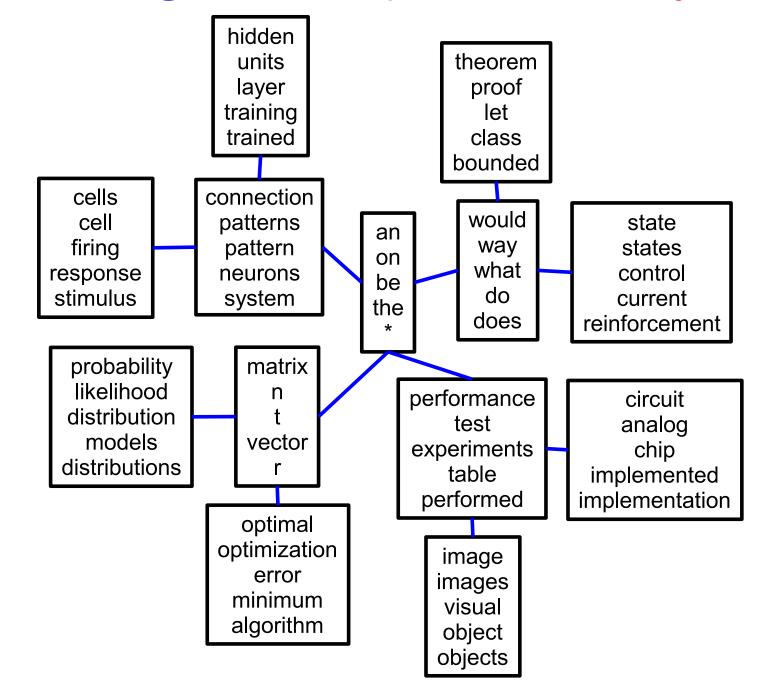


Variable selected after its ancestors (Zhao et al., 2009; Bach, 2008)

Hierarchical dictionary learning Modelling of text corpora

- Each document is modelled through word counts
- Low-rank matrix factorization of word-document matrix
- Probabilistic topic models (Blei et al., 2003)
 - Similar structures based on non parametric Bayesian methods (Blei et al., 2004)
 - Can we achieve similar performance with simple matrix factorization formulation?

Modelling of text corpora - Dictionary tree



Submodular functions and structured sparsity Examples

- From $\Omega(w)$ to F(A): provides new insights into existing norms
 - Grouped norms with **overlapping** groups (Jenatton et al., 2009a)

$$\Omega(w) = \sum_{G \in \mathbf{H}} \|w_G\|_{\infty} \Rightarrow F(A) = \operatorname{Card}(\{G \in \mathbf{H}, G \cap A \neq \emptyset\})$$

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- Justification not only limited to allowed sparsity patterns
- From F(A) to $\Omega(w)$: provides new sparsity-inducing norms
 - $-F(A) = g(\operatorname{Card}(A)) \Rightarrow \Omega$ is a combination of **order statistics**
 - Non-factorial priors for supervised learning: Ω depends on the eigenvalues of $X_A^{\top}X_A$ and not simply on the cardinality of A

Unified optimization algorithms

- Polyhedral norm with $O(3^p)$ faces and extreme points
 - Not suitable to linear programming toolboxes
- Subgradient ($w \mapsto \Omega(w)$ non-differentiable)
 - subgradient may be obtained in polynomial time ⇒ too slow

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 - subgradient may be obtained in polynomial time ⇒ too slow
- Proximal methods (e.g., Beck and Teboulle, 2009)
 - $-\min_{w\in\mathbb{R}^p} L(y,Xw) + \lambda\Omega(w)$: differentiable + non-differentiable
 - Efficient when (P): $\min_{w \in \mathbb{R}^p} \frac{1}{2} ||w v||_2^2 + \lambda \Omega(w)$ is "easy"
 - Fact: (P) is equivalent to submodular function minimization

Optimization for sparsity-inducing norms (see Bach, Jenatton, Mairal, and Obozinski, 2011)

Gradient descent as a proximal method (differentiable functions)

$$- w_{t+1} = \arg\min_{w \in \mathbb{R}^p} L(w_t) + (w - w_t)^{\top} \nabla L(w_t) + \frac{B}{2} ||w - w_t||_2^2$$
$$- w_{t+1} = w_t - \frac{1}{B} \nabla L(w_t)$$

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$$ullet$$
 Problems of the form: $\min_{w\in\mathbb{R}^p}L(w)+\lambda\Omega(w)$

$$- w_{t+1} = \arg\min_{w \in \mathbb{R}^p} L(w_t) + (w - w_t)^{\top} \nabla L(w_t) + \lambda \Omega(w) + \frac{B}{2} ||w - w_t||_2^2$$

- $-\Omega(w) = ||w||_1 \Rightarrow$ Thresholded gradient descent
- Similar convergence rates than smooth optimization
 - Acceleration methods (Nesterov, 2007; Beck and Teboulle, 2009)

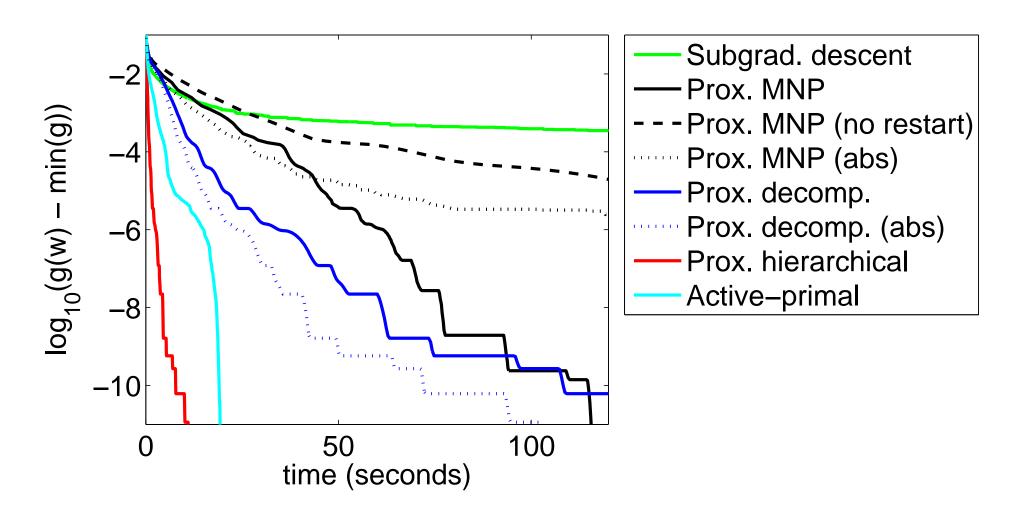
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 - Fact: (P) is equivalent to submodular function minimization

Active-set methods

Comparison of optimization algorithms

- Tree-based regularization (p = 511)
- See Bach et al. (2011) for larger-scale problems



Unified theoretical analysis

Decomposability

- Key to theoretical analysis (Negahban et al., 2009)
- **Property**: $\forall w \in \mathbb{R}^p$, and $\forall J \subset V$, if $\min_{j \in J} |w_j| \geqslant \max_{j \in J^c} |w_j|$, then $\Omega(w) = \Omega_J(w_J) + \Omega^J(w_{J^c})$

Support recovery

Extension of known sufficient condition (Zhao and Yu, 2006;
 Negahban and Wainwright, 2008)

High-dimensional inference

- Extension of known sufficient condition (Bickel et al., 2009)
- Matches with analysis of Negahban et al. (2009) for common cases

Support recovery - $\min_{w \in \mathbb{R}^p} \frac{1}{2n} ||y - Xw||_2^2 + \lambda \Omega(w)$

Notation

$$-\rho(J) = \min_{B \subset J^c} \frac{F(B \cup J) - F(J)}{F(B)} \in (0,1]$$
 (for J stable)

$$-c(J) = \sup_{w \in \mathbb{R}^p} \Omega_J(w_J) / ||w_J||_2 \leqslant |J|^{1/2} \max_{k \in V} F(\{k\})$$

Proposition

- Assume $y = Xw^* + \sigma\varepsilon$, with $\varepsilon \sim \mathcal{N}(0, I)$
- J= smallest stable set containing the support of w^*
- Assume $\nu = \min_{j, w_j^* \neq 0} |w_j^*| > 0$
- Let $Q = \frac{1}{n} X^{\top} X \in \mathbb{R}^{p \times p}$. Assume $\kappa = \lambda_{\min}(Q_{JJ}) > 0$
- Assume that for $\eta > 0$, $\left| (\Omega^J)^* [(\Omega_J(Q_{JJ}^{-1}Q_{Jj}))_{j \in J^c}] \right| \leq 1 \eta$
- If $\lambda \leqslant \frac{\kappa \nu}{2c(J)}$, \hat{w} has support equal to J, with probability larger than $1-3P\big(\Omega^*(z)>\frac{\lambda\eta\rho(J)\sqrt{n}}{2\sigma}\big)$
- -z is a multivariate normal with covariance matrix Q

Consistency - $\min_{w \in \mathbb{R}^p} \frac{1}{2n} ||y - Xw||_2^2 + \lambda \Omega(w)$

Proposition

- Assume $y = Xw^* + \sigma\varepsilon$, with $\varepsilon \sim \mathcal{N}(0, I)$
- -J = smallest stable set containing the support of w^*
- Let $Q = \frac{1}{n} X^{\top} X \in \mathbb{R}^{p \times p}$.
- $\text{ Assume that } \forall \Delta \text{ s.t. } \Omega^J(\Delta_{J^c}) \leqslant 3\Omega_J(\Delta_J), \ \Delta^\top Q \Delta \geqslant \kappa \|\Delta_J\|_2^2 \\ \text{ Then } \left[\Omega(\hat{w}-w^*) \leqslant \frac{24c(J)^2\lambda}{\kappa\rho(J)^2}\right] \text{ and } \left[\frac{1}{n}\|X\hat{w}-Xw^*\|_2^2 \leqslant \frac{36c(J)^2\lambda^2}{\kappa\rho(J)^2}\right]$

with probability larger than $1 - P(\Omega^*(z) > \frac{\lambda \rho(J)\sqrt{n}}{2\sigma})$

- -z is a multivariate normal with covariance matrix Q
- Concentration inequality (z normal with covariance matrix Q):
 - $-\mathcal{T}$ set of stable inseparable sets
 - Then $P(\Omega^*(z) > t) \leqslant \sum_{A \in \mathcal{T}} 2^{|A|} \exp\left(-\frac{t^2 F(A)^2/2}{1^T \Omega_{AA} 1}\right)$

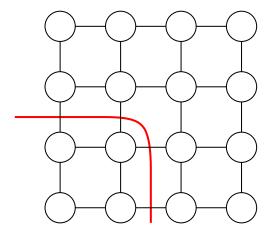
Symmetric submodular functions (Bach, 2011)

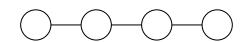
- ullet Let $F:2^V o \mathbb{R}$ be a symmetric submodular set-function
- Proposition: The Lovász extension f(w) is the convex envelope of the function $w \mapsto \max_{\alpha \in \mathbb{R}} F(\{w \geqslant \alpha\})$ on the set $[0,1]^p + \mathbb{R}1_V = \{w \in \mathbb{R}^p, \max_{k \in V} w_k \min_{k \in V} w_k \leqslant 1\}.$
- Shaping all level sets

Symmetric submodular functions - Examples

- From $\Omega(w)$ to F(A): provides new insights into existing norms
 - Cuts total variation

$$F(A) = \sum_{k \in A, j \in V \setminus A} d(k, j) \quad \Rightarrow \quad f(w) = \sum_{k, j \in V} d(k, j) (w_k - w_j)_+$$



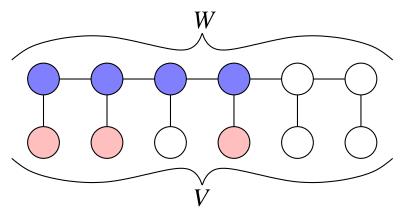


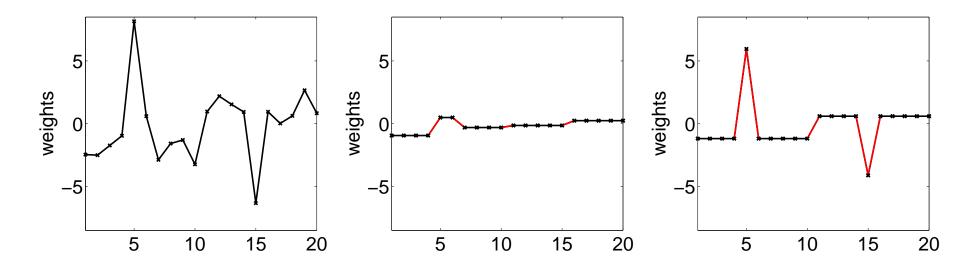
- NB: graph may be directed
- Application to change-point detection (Tibshirani et al., 2005;
 Harchaoui and Lévy-Leduc, 2008)

Symmetric submodular functions - Examples

- From F(A) to $\Omega(w)$: provides new sparsity-inducing norms
 - Regular functions (Boykov et al., 2001; Chambolle and Darbon, 2009)

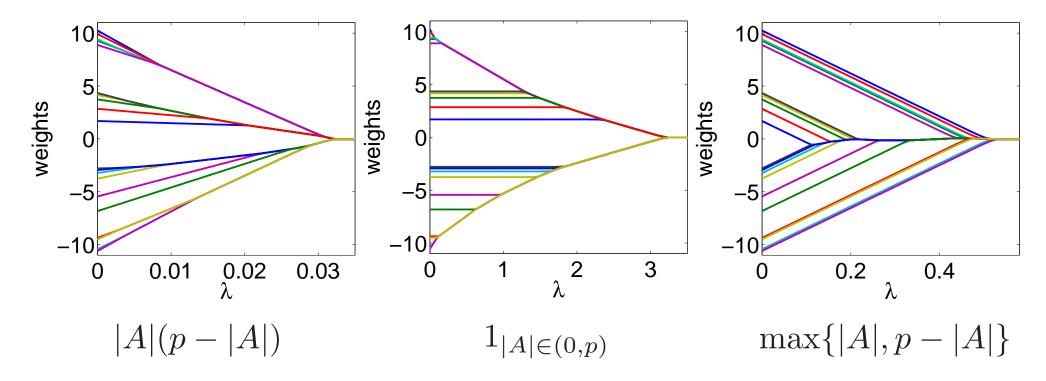
$$F(A) = \min_{B \subset W} \sum_{k \in B, \ j \in W \setminus B} d(k, j) + \lambda |A \Delta B|$$





Symmetric submodular functions - Examples

- From F(A) to $\Omega(w)$: provides new sparsity-inducing norms
 - $-F(A) = g(\operatorname{Card}(A)) \Rightarrow$ priors on the size and numbers of clusters



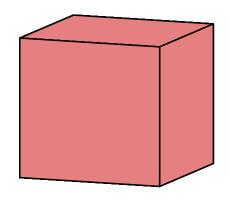
 Convex formulations for clustering (Hocking, Joulin, Bach, and Vert, 2011)

ℓ_2 -relaxation of combinatorial penalties (Obozinski and Bach, 2012)

- Main result of Bach (2010):
 - -f(|w|) is the convex envelope of $F(\operatorname{Supp}(w))$ on $[-1,1]^p$

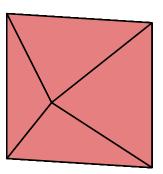
Problems:

- Limited to submodular functions
- Limited to ℓ_{∞} -relaxation: undesired artefacts



$$F(A) = \min\{|A|, 1\}$$

$$\Omega(w) = ||w||_{\infty}$$



$$F(A) = 1_{\{A \cap \{1\} \neq \emptyset\}} + 1_{\{A \cap \{2,3\} \neq \emptyset\}}$$

$$\Omega(w) = |w_1| + ||w_{\{2,3\}}||_{\infty}$$

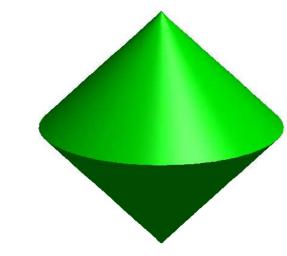
ℓ_2 -relaxation of submodular penalties (Obozinski and Bach, 2012)

ullet F a nondecreasing submodular function with Lovász extension f

• Define
$$\Omega_2(w) = \min_{\eta \in \mathbb{R}^p_+} \frac{1}{2} \sum_{i \in V} \frac{|w_i|^2}{\eta_i} + \frac{1}{2} f(\eta)$$

- NB: general formulation (Micchelli et al., 2011; Bach et al., 2011)
- **Proposition 1**: Ω_2 is the convex envelope of $w \mapsto F(\operatorname{Supp}(w)) \|w\|_2$
- Proposition 2: Ω_2 is the *homogeneous* convex envelope of $w\mapsto \frac{1}{2}F(\operatorname{Supp}(w))+\frac{1}{2}\|w\|_2^2$
- Jointly penalizing and regularizing
 - Extension possible to ℓ_q , q > 1

From ℓ_{∞} to ℓ_2 Removal of undesired artefacts



$$F(A) = 1_{\{A \cap \{3\} \neq \emptyset\}} + 1_{\{A \cap \{1,2\} \neq \emptyset\}}$$
$$\Omega_2(w) = |w_3| + ||w_{\{1,2\}}||_2$$



$$F(A) = 1_{\{A \cap \{1,2,3\} \neq \emptyset\}} + 1_{\{A \cap \{2,3\} \neq \emptyset\}} + 1_{\{A \cap \{2\} \neq \emptyset\}}$$

ullet Extension to non-submodular functions + tightness study: see Obozinski and Bach (2012)

Beyond submodular functions?

- \bullet Let F be any set-function
- "Edmonds extension": homogeneous convex envelope of $w \mapsto F(\operatorname{Supp}(w))$ on $[0,1]^p$ equal to

$$f(w) = \sup_{\forall A \subseteq V, \ s(A) \leqslant F(A)} w^{\top} s = \sup_{s \in P(F)} w^{\top} s$$

- When is it an extension of F?
- Lower combinatorial envelope: $G(B) = f(1_B) = \sup_{s \in P(F)} s(B)$
 - $-G \leqslant F$
 - Property: idempotent operation
- A new class of set-functions: functions for which G = F

Conclusion

- Structured sparsity for machine learning and statistics
 - Many applications (image, audio, text, etc.)
 - May be achieved through structured sparsity-inducing norms
 - Link with submodular functions: unified analysis and algorithms

Submodular functions to encode discrete structures

Conclusion

• Structured sparsity for machine learning and statistics

- Many applications (image, audio, text, etc.)
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- Link with submodular functions: unified analysis and algorithms
 Submodular functions to encode discrete structures

On-going work on machine learning and submodularity

- Submodular function maximization
- Importing concepts from machine learning (e.g., graphical models)
- Multi-way partitions for computer vision
- Online learning

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