

### 3.3. Chop-and-Expand Context-Free Parser in a logic language

Recall that in logic programming  $[A|X]$  represents the list whose head is the element  $A$  and whose tail is the list  $X$ .

Here is a logic program which realizes a Chop-and-Expand parser:

```

1.  parse( $G, [], []$ )  $\leftarrow$ 
2.  parse( $G, [A|X], [A|Y]$ )  $\leftarrow$  terminal( $A$ ), parse( $G, X, Y$ )           // CHOP
3.  parse( $G, [A|X], Y$ )  $\leftarrow$  nonterminal( $A$ ), member( $A \rightarrow B, G$ ), // EXPAND
      append( $B, X, Z$ ), parse( $G, Z, Y$ )
4.  member( $A, [A|X]$ )  $\leftarrow$ 
5.  member( $A, [B|X]$ )  $\leftarrow$  member( $A, X$ )
6.  append( $[], L, L$ )  $\leftarrow$ 
7.  append( $[A|X], Y, [A|Z]$ )  $\leftarrow$  append( $X, Y, Z$ )

```

together with the clauses which define which are the terminal symbols and the non-terminal symbols.

The first argument of *parse* is a context-free grammar, the second argument is a list of terminal symbols and nonterminal symbols (that is, a sentential form), and the third argument is a word represented as a list of terminal symbols. We assume that context-free grammars are represented as lists of productions of the form  $x \rightarrow y$ , where  $x$  is a nonterminal symbol and  $y$  is a list of terminal and nonterminal symbols.

We have that  $\textit{parse}(G, [s], W)$  holds iff from the symbol  $s$  we can derive the word  $W$  using the grammar  $G$ .

EXAMPLE 3.3.1. The grammar  $G = \langle \{a, b\}, \{S\}, \{S \rightarrow aSb, S \rightarrow ab\}, S \rangle$  is represented by the clauses:

```

terminal( $a$ )  $\leftarrow$ 
terminal( $b$ )  $\leftarrow$ 
nonterminal( $s$ )  $\leftarrow$ 

```

together with the list  $[s \rightarrow [a, s, b], s \rightarrow [a, b]]$  which represents its productions. The left hand side of the first production is assumed to be the start symbol. For this grammar  $G$  the goal  $\leftarrow \textit{parse}([s \rightarrow [a, s, b], s \rightarrow [a, b]], [s], [a, a, b, b])$  is true.

Note that in the clause  $\textit{member}(A, [B|X]) \leftarrow \textit{member}(A, X)$ , we do not require that  $A$  is different from  $B$ . Indeed with this clause, the query of the form  $\textit{member}(A, l)$ , where  $A$  is a variable and  $l$  is a ground list, generates by backtracking all members of the list  $l$  with all their multiplicity.  $\square$